Imperial College London

### Temporal blocking of finite-difference stencil operators with sparse off-the-grid sources

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### Talk outline



- Motivation: speed up computationally expensive scientific simulations involving the solution of PDEs modelling wave equations through explicit FD methods (seismic and medical imaging)
- Accelerating through cache optimizations, more specifically through temporal blocking
- Enabling temporal blocking on practical wave-propagation simulations is complicated as they consist of sparse
   "off-the-grid" operators (Not a typical stencil benchmark!)
   => Applicability issues
- We present **an approach to overcome limitations** and enable TB
- Experimental results show improved performance





### Modelling practical applications

- Stencils everywhere, not only though. What else?
- Remarkable amount of work in the past on optimizing stencils... (Parallelism, cache optimizations, accelerators)
- Sources injecting and receivers interpolating at sparse off-the-grid coordinates.
   Non-conventional update patterns.
- Usually their coordinates are not aligned with the computational grid. How do we iterate over them?







A 1d 3pt stencil update

A 3d-19pt stencil update



Off-the-grid operators (Source injection/Receiver interpolation)

### Sparse off-the-grid operators





How a seismic survey looks like

Source: KrisEnergy 2021

### Sparse off-the-grid operators



- How a seismic survey looks like
- Discretizing the computational domain (the FD-grid). Solution computed on the points

Source: KrisEnergy 2021



### Sparse off-the-grid operators



- How a seismic survey looks
   like
- Discretizing the computational domain (the FD-grid). Solution computed on the points
- Not-aligned "off-the-grid" operators exist (source injection/receiver interpolation)

Source: KrisEnergy 2021



### A typical time-stepping loop with source injection

- Stencil computation, points on the domain are updated as a weighted function of neighbouring points
- Then we update the points that are affected from sources
- Each source has 3D coordinates
- Indirect accesses are used to scatter injection to neighbouring points
- Sources are aligned in time, same time index with points
- Sources are NOT aligned in space, off-the grid discretization

**Listing 1:** A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.

1 <b>f</b>	or $t = 1$ to nt do
2	for $x = 1$ to nx do
3	<b>for</b> $y = 1$ to ny <b>do</b>
4	for $z = 1$ to nz do
5	$        A(t, x, y, z) \equiv u[t, x, y, z] = u[t-1, x, y, z] + \sum_{r=1}^{r=so/2} w_r ($
	u[t-1, x - r, y, z] + u[t-1, x + r, y, z] + u[t-1, x, y - r, z] +
	$\left  \begin{array}{c} u[t-1, x, y + r, z] + u[t-1, x, y, z - r] + u[t-1, x, y, z + r] \end{array} \right ;$
6	foreach s in sources do // For every source
7	<b>for</b> $i = 1$ to np do // Get the points affected
8	xs, ys, zs = map(s, i) // through indirection
9	u[t, xs, ys, zs] + = f(src(t, s)) // add their impact
	on the field





### Applying loop-blocking

Loop blocking (aka space blocking, loop tiling):

- Decompose grids into blocks/tiles. Iteration space partitioned to smaller chunks/blocks
- Improved data locality  $\Rightarrow$  Increased performance (Rich literature)
- Sparse off-the-grid operators are iterated as without blocking





### Applying temporal-blocking

Temporal blocking (Time-Tiling):

- Space blocking but data reuse is extended to time-dimension.
- Update grid points in future where (space) and when (time) possible
- Rich literature, several variants of temporal blocking, shapes, schemes
  - -Wave-front / Skewed (Approach followed in the paper)
  - -Diamonds, Trapezoids, Overlapped, Hybrid models





### Off-the-grid operators: the issue



- Data dependences violations happen while a temporal update
- Source injection is in a different iteration space
- When a sparse operator exists in the boundary between space-time blocks, the order of updates is not preserved
- Solution: Need to align off-the-grid operators



### Off-the-grid operators: the issue



- Data dependences violations happen while a temporal update
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- Solution: Need to align off-the-grid operators





## Methodology

- A negligible-cost scheme to precompute the source injection contribution.
- Align source injection data dependences to the grid
- This scheme is applicable to other fields as well (e.g. medical imaging)

### Iterate over sources and store indices of affected points



- Inject to a **zero-valued initialized grid** for one (or a few more) timesteps
- **Hypothesis:** non-zero source-injection values at the first time-steps
- Independent of the injection and interpolation type (e.g. non-linear injection)

**Listing 2:** Source injection over an empty grid. No PDE stencil update is happening.

- 1 for t = 1 to 2 do 2 foreach s in sources do 3 for i = 1 to np do 4 |xs, ys, zs = map(s, i);5 |u[t, xs, ys, zs] + = f(src(t, s))
  - Then, we **store the non-zero** grid point coordinates

Generate sparse binary mask, unique IDs and decompose wavefields



- Perform source injection to decompose • the off-the-grid wavefields to on-the-grid per point wavefields.
- Inject sources to sources

	Off-the-grid	Aligned
len(sources)	n_src	n_aff_pts
len(sources.coords)	(n_src, 3)	(n_aff_pts, 3)
len(sources.data)	(n_src, nt)	(n_aff_pts, nt)

Listing 3: Decomposing the source injection wavefields.

for t = 1 to nt do 

**foreach** s *in* sources **do** 

**for** 
$$i = 1$$
 to np do

$$|xs, ys, zs = map(s, i);$$

$$s \quad | \quad | \operatorname{src\_dcmp}[t, \operatorname{SID}[xs, ys, zs]] + = f(\operatorname{src}(t, s);$$



(a) Sources are sparsely distributed at off-the-grid posi-

o npts

tions.

Y



Х

(b) Identify unique points affected (SM).

	-							
	•	0	0	0		•	0	0
	0	0	ο	0	0	0	ο	0
	0	ο	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
Y	0	0	ο	0	0	0	o	0
	0		•	0	0	0	0	0
	0	•		o	0	0	0	
	0	0	0	0	0	0	0	•
	L				v			
					~			

(c) Assign a unique ID to every affected point (SID).

X

(d) Sources are aligned with grid positions.

### Fuse iteration spaces



- Indirection mapping has changed. We still use indirections but now they are on the point.
- By using the aligned structure, we fuse the source injection loop inside the kernel update iteration space.
- The source mask SM is used to add (if 1) or not (if 0) the impact and SID is used to indirect to the impact values using the traversed grid coordinates.

Listing 4: Stencil kernel update with fused source injection.

1	for <i>t</i> = 1 <i>to</i> nt do	
2	for $x = 1$ to nx do	
3	for y = 1 to ny do	
4	<b>for</b> $z = 1$ to nz do	
5	A(t, x, y, z, s); SIMD2 (A)(X512)	
6	for $z^2 = 1$ to $nz$ do SIVID? (AVASIZ)	
7	$        u[t, x, y, z2] + = SM[x, y, z2] * src_dcmp[t, SID[x, y, z2]$	]];



### Reducing the iteration space size

- Lots of redundant ops due to sparsity
- A schedule to perform only necessary operations
- Aggregate NZ along the z-axis keeping count of them in a reduced-size structure named *nnz\_mask*
- Reduce the size of SID by cutting off zero z-slices



**Listing 5:** Stencil kernel update with reduced size iteration space for source injection.

1	for $t = 1$ to nt do
2	for $x = 1$ to nx do
3	<b>for</b> y = 1 to ny <b>do</b>
4	<b>for</b> $z = 1$ to nz do
5	A(t, x, y, z, s);
6	for $z^2 = 1$ to nnz_mask[x][y] do
7	$      I(t, x, y, z, s) \equiv \{ \text{ zind} = \text{Sp}SID[x, y, z2]; $
8	$    u[t, x, y, z2] += \operatorname{src\_dcmp}[t, SID[x, y, zind]]; \}$

**Listing 1:** A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.

1 1	for $t = 1$ to nt do
2	for $x = 1$ to nx do Non-aligned
3	10r  y = 1  10  ny  d0
4	for $z = 1$ to nz do
5	$A(t, x, y, z) \equiv u[t, x, y, z] = u[t-1, x, y, z] + \sum_{r=1}^{r=so/2} w_r \Big($
	u[t-1, x - r, y, z] + u[t-1, x + r, y, z] + u[t-1, x, y - r, z] +
	u[t-1, x, y + r, z] + u[t-1, x, y, z - r] + u[t-1, x, y, z + r]);
6	foreach s in sources do // For every source
7	<b>for</b> $i = 1$ to np do // Get the points affected
8	$x_s, y_s, z_s = map(s, i) // through indirection$
9	u[t, xs, ys, zs] + = f(src(t, s)) / / add their impact
	on the field

**Listing 5:** Stencil kernel update with fused - reduced size iteration space - source injection.

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for 
$$t = 1$$
 to nt do  
for  $x = 1$  to nx do  
for  $y = 1$  to ny do  
for  $z = 1$  to nz do  
 $|A(t, x, y, z, s);$   
for  $z^2 = 1$  to nnz\_mask[x][y] do  
 $|I(t, x, y, z, s) \equiv \{ \text{ zind = Sp_SM[x, y, z2];} \\ u[t, x, y, zd] += \{ \text{SM[x, y, zind]} \}$ 





Everything so far, automated in Devito DSL



### Applying wave-front temporal blocking



- TB with manually editing the Devito generated code
- Skewing factor depends on data dependency distances (higher for higher SO, multigrid)



Figure from YASK, Yount et. al (2016)



(a) The figure shows multiple wave-front tiles evaluated sequentially, partially adapted from [15].



(b) The figure shows multiple wave-front tiles evaluated sequentially in multigrid stencil codes. **Listing 6:** The figure shows the loop structure after applying our proposed scheme.



```
for t tile in time tiles do
  for xtile in xtiles do
                                              Iterate space-time tiles
     for ytile in ytiles do
        for t in t_tile do
                                              Time-stepping in the tile and loop-blocking within
          OpenMP parallelism
                                              the tile. Collapse outer loops that are loop-blocked
           for xblk in xtile do
             for yblk in ytile do
                for x in xblk do
                   for y in yblk do
                                               No loop blocking on z-dim, full stride for
                      SIMD vectorization
                                               max-vectorization performance
                      for z = 1 to nz do
                        |A(t, x - time, y - time, z, s);
                      for z<sup>2</sup> = 1 to nnz_mask[x][y] do
                        |I(t, x - time, y - time, z2, s);
```



### Experimental evaluation: the models

### Isotropic Acoustic

Generally known, single scalar PDE, laplacian like, low cost

### • Isotropic Elastic

Coupled system of a vectorial and tensorial PDE, explosive source, increased data movement, first order in time, cross-loop data dependencies

• Anisotropic Acoustic (aka TTI)

Industrial applications, rotated laplacian, coupled system of two scalar PDEs

Industrial-level, 512^3 grid points, 512ms simulation time, damping fields ABCs



Velocity field, TTI wave propagation after 512ms

### Experimental evaluation: the results





Azure model Architecture	E16s v3 Broadwell	E32s v3 Skylake
vCPUs	16	32
GiB memory	128	256
Model name	E5-2673 v4	8171M
CPUs	16	32
Thread(s) per core	2	2
Core(s) per socket	8	16
Socket(s)	1	1
NUMA node(s)	1	1
Model	79	85
CPU MHz	2300	2100
L1d cache	32K	32K
L1i cache	32K	32K
L2 cache	256K	1024K
L3 cache	51200K	36608K



(b) Throughput speed-up of kernels for Skylake.

- Benchmark on Azure VMs
- GCC, ICC
- Thread pinning
- OpenMP, SIMD
- Aggressive auto-tuning



- Kernels are flop-optimized through Devito.
- Gpts/s aka Gcells/s: time to solution metric in stencil computations
- (!) High Gflops/s do not guarantee a faster solution.

TABLE I: VM specification

### Cache-aware roofline model





Space order: •∆4 •O8 •□12 Temporal Blocking Spatial Blocking

Broadwell, isotropic acoustic, 512<sup>3</sup> grid points, 512ms

### Conclusions



- We presented an approach to apply temporal blocking to stencil kernels with sparse off-the-grid operators.
- The additional cost is negligible compared to the achieved gains.
- Solution built on top of Devito-DSL
- Performance gains of up to 1.6x on low order (4) and 1.15x on medium order (8).

 Integration to DSL Current WIP . User will get out-of the box time tiled code for all PDEs!

Future plans • MPI-aware scheme

- GPUs
- High-order stencils





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### Corner cases, increasing number of sources



### The generated C code - stencil update

```
#pragma omp tor collapse(1) schedule(dynamic,1)
             for (int x0 blk0 = x m; x0 blk0 <= x M; x0 blk0 += x0 blk0 size)</pre>
                   for (int y0 blk0 = y m; y0 blk0 <= y M; y0 blk0 += y0 blk0 size)
                          for (int x = x0 blk0; x <= x0 blk0 + x0 blk0 size - 1; x += 1)
                                for (int y = y0_blk0; y <= y0_blk0 + y0_blk0_size - 1; y += 1)</pre>
                                       #pragma omp simd aligned(damp,uref,vp:32)
                                       for (int z = z m; z <= z M; z += 1)
                                             float r14 = -2.84722222F*uref[t1][x + 8][y + 8][z + 8];
                                             float r13 = 1.0/dt:
                                             float r12 = 1.0/(dt*dt);
                                             float r11 = 1.0/(vp[x + 8][y + 8][z + 8]*vp[x + 8][y + 8][z + 8]);
                                             uref[t0][x + 8][y + 8][z + 8] = (r11*(-r12*(-2.0F*uref[t1][x + 8][y + 8][z + 8] + 8)]
uref[t_2][x + 8][y + 8][z + 8]) + r_{13}(damp[x + 1][y + 1][z + 1]*uref[t_1][x + 8][y + 8][z + 8])
(r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 4] + uref[t1][x + 8][y + 8][z + 12]) +
2.53968254e-2F*(uref[t1][x + 8][y + 8][z + 5] + uref[t1][x + 8][y + 8][z + 11]) -
2.0e-1F*(uref[t1][x + 8][y + 8][z + 6] + uref[t1][x + 8][y + 8][z + 10]) + 1.6F*(uref[t1][x + 8][y + 8][z + 10]) + 1.6F*(uref[t1][x + 8][y + 8][y + 8][z + 10]) + 1.6F*(uref[t1][x + 8][y + 8][y + 8][z + 10]) + 1.6F*(uref[t1][x + 8][y + 8][
[y + 8][z + 7] + uref[t1][x + 8][y + 8][z + 9])/((h z*h z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h z*h z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h z*h z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h z*h z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h z*h z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][y + 8][y + 8][y + 8][y + 8](y + 8][y + 8](y + 8][y + 8][y
 + 8][y + 4][z + 8] + uref[t1][x + 8][y + 12][z + 8]) + 2.53968254e-2F*(uref[t1][x + 8][y + 5][z + 8])
B] + uref[t1][x + 8][y + 11][z + 8]) - 2.0e-1F*(uref[t1][x + 8][y + 6][z + 8] + uref[t1][x + 8][
 + 10[z + 8] + 1.6F*(uref[t1][x + 8][y + 7][z + 8] + uref[t1][x + 8][y + 9][z + 8]))/((h_y*h_y))
+ (r14 - 1.78571429e - 3F*(uref[t1][x + 4][y + 8][z + 8] + uref[t1][x + 12][y + 8][z + 8]) +
 2.53968254e-2F*(uref[t1][x + 5][y + 8][z + 8] + uref[t1][x + 11][y + 8][z + 8]) -
2.0e-1F*(uref[t1][x + 6][y + 8][z + 8] + uref[t1][x + 10][y + 8][z + 8]) + 1.6F*(uref[t1][x + 7])
[y + 8][z + 8] + uref[t1][x + 9][y + 8][z + 8])/((h x*h x))/(r11*r12 + r13*damp[x + 1][y + 1][
 + 11);
                                      }
      }
}
     }
```

### The generated C code - source injection



```
/* Begin section1 */
    #pragma omp parallel num threads(nthreads nonaffine)
      int chunk size = (int)(fmax(1, (1.0F/3.0F)*(p src M - p src m + 1)/nthreads nonaffine));
     #pragma omp for collapse(1) schedule(dynamic,chunk size)
     for (int p src = p src m; p src <= p src M; p src += 1)
       int ii src 0 = (int)(floor((-o x + src coords[p src][0])/h x));
       int ii src 1 = (int)(floor((-o y + src coords[p src][1])/h y));
       int ii src 2 = (int)(floor((-o z + src coords[p src][2])/h z));
       int ii src 3 = (int)(floor((-o z + src coords[p src][2])/h z)) + 1;
       int ii src 4 = (int)(floor((-o y + src coords[p src][1])/h y)) + 1;
       int ii src 5 = (int)(floor((-o x + src coords[p src][0])/h x)) + 1;
       float px = (float)(-h x*(int)(floor((-o x + src coords[p src][0])/h x)) - o x + src coords[p src][0]);
       float pv = (float)(-h v*(int)(floor((-o y + src coords[p src][1])/h_y)) - o_y + src_coords[p_src][1]);
       float pz = (float)(-h z*(int)(floor((-o z + src coords[p src][2])/h z)) - o z + src coords[p src][2]);
       if (ii src 0 >= x m - 1 && ii src 1 >= y m - 1 && ii src 2 >= z m - 1 && ii src 0 <= x M + 1 && ii src 1
<= v M + 1 && ii src 2 <= z M + 1)
          float r0 = 4.49016082216644F*(vp[ii_src_0 + 8][ii_src_1 + 8][ii_src_2 + 8]*vp[ii_src_0 + 8][ii_src_1 + 8]
[ii src 2 + 8])*(-px*py*pz/(h_x*h_y*h_z) + px*py/(h_x*h_y) + px*pz/(h_x*h_z) - px/h_x + py*pz/(h_y*h_z) - py/h_y -
pz/h z + 1)*src[time][p src];
          #pragma omp atomic update
         uref[t0][ii src 0 + 8][ii src 1 + 8][ii src 2 + 8] += r0;
       if (ii src 0 >= x m - 1 && ii src 1 >= y m - 1 && ii src 3 >= z m - 1 && ii src 0 <= x M + 1 && ii src 1
<= y M + 1 && ii src 3 <= z M + 1)
       ſ
          float r1 = 4.49016082216644F*(vp[ii src 0 + 8][ii src 1 + 8][ii src 3 + 8]*vp[ii src 0 + 8][ii src 1 + 8]
[ii src 3 + 8])*(px*py*pz/(h x*h y*h z) - px*pz/(h x*h z) - py*pz/(h y*h z) + pz/h z)*src[time][p src];
          #pragma omp atomic update
         uref[t0][ii_src_0 + 8][ii_src_1 + 8][ii_src_3 + 8] += r1;
       if (ii src 0 >= x m - 1 && ii src 2 >= z m - 1 && ii src 4 >= y m - 1 && ii src 0 <= x M + 1 && ii src 2
<= z M + 1 && ii src 4 <= y M + 1)
        Ł
          float r2 = 4.49016082216644F*(vp[ii src 0 + 8][ii src 4 + 8][ii src 2 + 8]*vp[ii src 0 + 8][ii src 4 + 8]
[ii src 2 + 8])*(px*py*pz/(h x*h v*h z) - px*py/(h x*h v) - pv*pz/(h v*h z) + pv/h v)*src[time][p src]:
```



### Cache aware roofline model

From here: <u>https://crd.lbl.gov/departments/computer-science/par/research/roofline/introduction/</u>

Effects of Cache Behavior on Arithmetic Intensity

The Roofline model requires an estimate of total data movement. On cache-based architectures, the 3C's cache model highlights the fact that there can be more than simply compulsory data movement. Cache capacity and conflict misses can increase data movement and reduce arithmetic intensity. Similarly, superfluous cache write-allocations can result in a doubling of data movement. The vector initialization operation x[i]=0.0 demands one write allocate and one write back per cache line touched. The write allocate is superfluous as all elements of that cache line are to be overwritten. Unfortunately, the presence of hardware stream prefetchers can make it very difficult to quantify how much beyond compulsory data movement actually occurred.

Algorithm 3: Source injection pseudocode.



for t = l to nt do foreach s in sources do 2 Discover affected # Find on the grid coordinates 3 src x min = floor(src coords[s][0], ox) 4 points 5  $src_x_max = ceil(src_coords[s][0], ox)$ # Compute weights 6 Weights of impact 7  $px = f(src\_coords[s][0], ox)$ # Unrolled for 8 points  $(2^\circ, 3D \text{ case})$ 8 if src\_x\_min, ... in grid then 9 Unrolled loop for  $r0 = v(src_x_min, \dots src[t][s]);$ 10  $u[t, src_x_{min}, ...] + = r0)$ each affected 11 point, compute 12 if src\_x\_max, ... in grid then injection part and  $r7 = v(src_x_max, \dots src[t][s]);$ 13 add to field  $u[t, src_x_max, ...] + = r7)$ 14 30













Algorithm 1: A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.

1 for t = 1 to nt do for x = 1 to  $n \times do$ 2 Non-aligned for y = 1 to ny do 3 for z = 1 to nz do 4  $\begin{vmatrix} A(t, x, y, z) \equiv u[t, x, y, z] = u[t-1, x, y, z] + \sum_{r=1}^{r=so/2} w_r \begin{bmatrix} u[t-1, x - r, y, z] + u[t-1, x + r, y, z] + u[t-1, x, y - r, z] + u[t-1, x, y - r, z] \end{bmatrix}$ 5 u[t-1, x, y + r, z] + u[t-1, x, y, z - r] + u[t-1, x, y, z + r]; foreach s in sources do 6 for *i* = 1 *to* np do 7 xs, ys, zs = map(s, i); 8 [u[t, xs, ys, zs] + = f(src(t, s))9

**Algorithm 6:** Stencil kernel update with fused - reduced size iteration space - source injection.

```
for t = 1 to nt do

for x = 1 to nx do

for y = 1 to ny do

for z = 1 to nz do

|A(t, x, y, z, s);

for z^2 = 1 to nnz_mask[x][y] do

|zind = Sp_SM[x, y, z];

u[t, x, y, z2] +=

|SM[x, y, zind] * src_dcmp[t, SID[x, y, zind]];
```

Aligned

**Listing 1:** A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.



**Listing 5:** Stencil kernel update with fused - reduced size iteration space - source injection.

for 
$$t = 1$$
 to nt do  
for  $x = 1$  to nx do  
for  $y = 1$  to ny do  
for  $z = 1$  to nz do  
 $|A(t, x, y, z, s);$   
for  $z^2 = 1$  to nnz\_mask[x][y] do  
 $|I(t, x, y, z, s) \equiv \{ \text{ zind } = \text{Sp}[SM[x, y, z^2]; u[t, x, y, z^2] += SM[x, y, zind] * \text{src_dcmp[t, SID[x, y, zind]]; } \}$ 
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Aligned to grid
Same OPS
Parallelism
SIMD (?)
Apply TB



### The transformation in Devito-DSL

u = TimeFunction(name="u", grid=model.grid, space\_order=so, time\_order=2)
src\_term = src.inject(field=u.forward, expr=src \* dt\*\*2 / model.m)
pde = model.m \* u.dt2 - u.laplace + model.damp \* u.dt
stencil = Eq(u.forward, solve(pde, u.forward))
op = Operator([stencil, src\_term])

### The transformation in Devito-DSL

# f : perform source injection on an empty grid

f = TimeFunction(name="f", grid=model.grid, space\_order=so, time\_order=2) src\_f
= src.inject(field=f.forward, expr=src \* dt\*\*2 / model.m)
op\_f = Operator([src\_f])
op\_f\_sum = op\_f.apply(time=3)

nzinds = np.nonzero(f.data[0]) # nzinds is a tuple

eqO = Eq(sp\_zi.symbolic\_max, nnz\_sp\_source\_mask[x, y] - 1, implicit\_dims=(time, x, y)) eq1 = Eq(zind, sp\_source\_mask[x, y, sp\_zi], implicit\_dims=(time, x, y, sp\_zi))

mask\_expr = source\_mask[x, y, zind] \* save\_src[time, source\_id[x, y, zind]]
eq2 = Inc(usol.forward[t+1, x, y, zind], mask\_expr, implicit\_dims=(time, x, y, sp\_zi)) pde\_2 =

model.m \* usol.dt2 - usol.laplace + model.damp \* usol.dt

stencil\_2 = Eq(usol.forward, solve(pde\_2, usol.forward))

### Fuse iteration spaces

- Indirection mapping has changed. We still use indirections but now they are on the point.
- By using the aligned structure, we fuse the source injection loop inside the kernel update iteration space.
- The source mask SM is used to add (if 1) or not (if 0) the impact and SID is used to indirect to the impact values using the traversed grid coordinates.

Listing 4:	Stencil kerne	el update	with	fused	source	injection.
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1	for $t = 1$ to nt do
2	for $x = 1$ to nx do
3	for y = 1 to ny do
4	<b>for</b> $z = 1$ to nz do
5	A(t, x, y, z, s);
6	for $z^2 = 1$ to nz do
7	$ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $