Imperial College London

TEMPORAL BLOCKING FOR WAVE-PROPAGATION KERNELS WITH SPARSE OFF-THE-GRID SOURCES

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Oil and Gas

High-Performance Computing Conference 2021

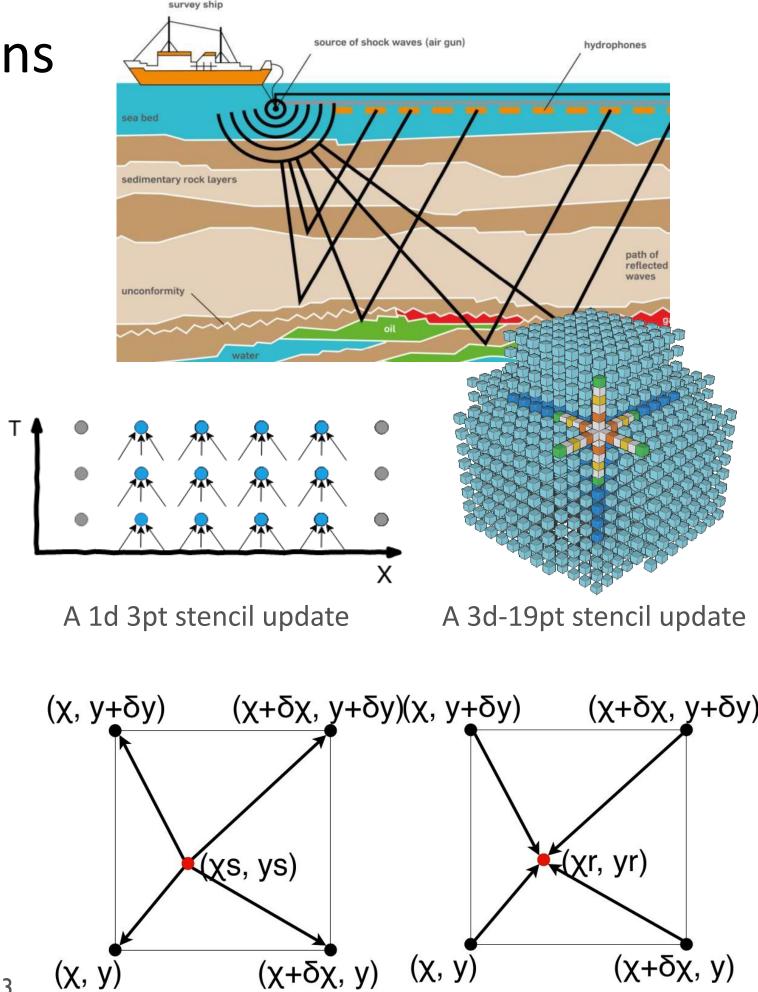
Rice University, Ken Kennedy Institute, Houston

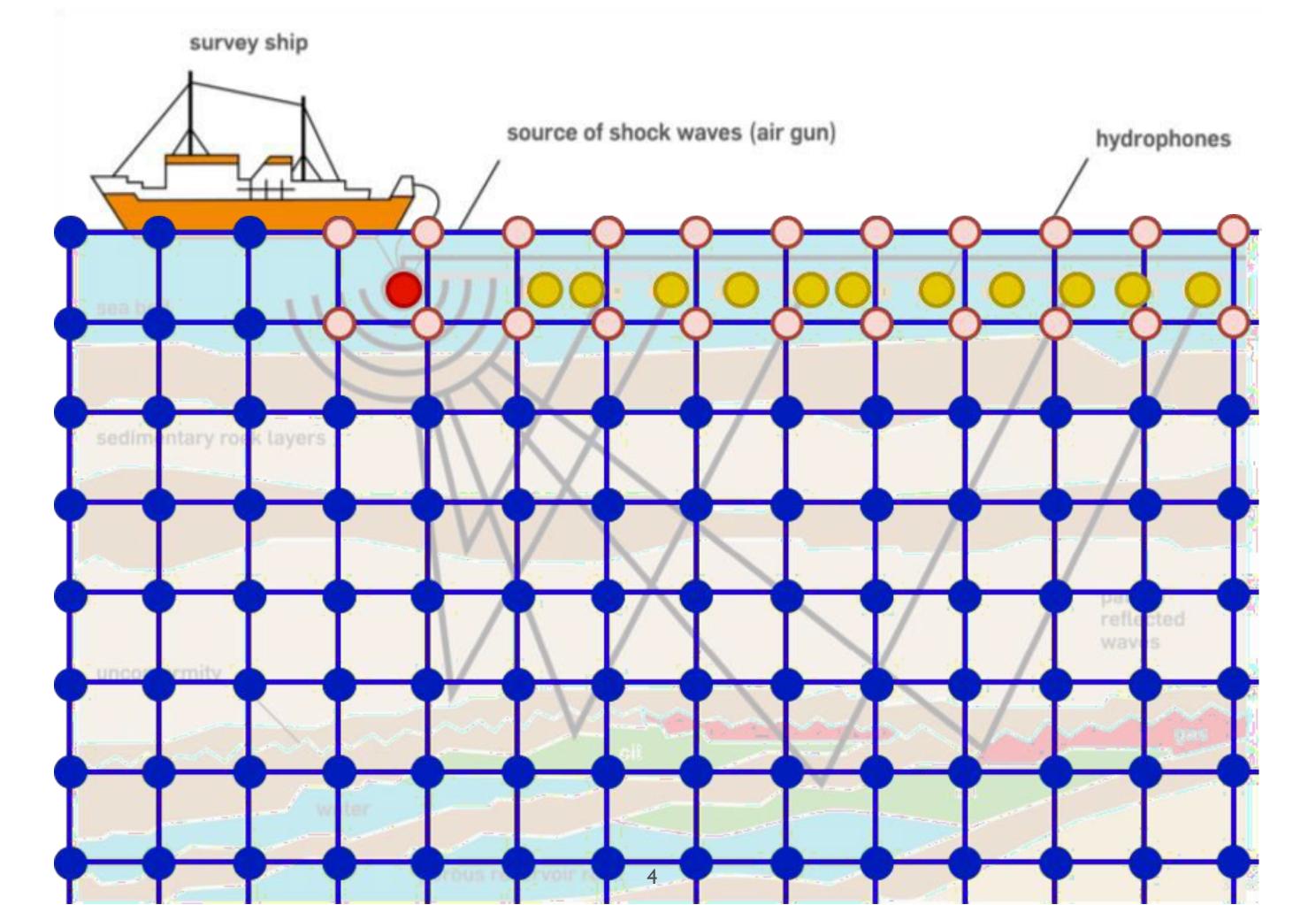
Motivation

- Accelerating wave-propagation kernels of practical interest through cache optimizations, more specifically through temporal blocking
- Enabling **Temporal Blocking** on practical wave-propagation simulations is complicated
- They consist of **sparse "off-the-grid"** operators (**Not** a typical stencil benchmark!)
- Applicability issues due to seismic imaging kernels consisting of sparse off-the-grid operators
- We present an approach to overcome limitations and enable TB
- Experimental results show improved performance

Modelling practical applications

- Stencils everywhere, but not only. What else?
- Vast literature on optimizing stencils... (Parallelism, cache optimizations, accelerators)
- Sources injecting and receivers interpolating at sparse off-the-grid coordinates. Non-conventional update patterns.
- Usually their coordinates are not aligned with the computational grid. How do we iterate over them?





A typical time-stepping loop with source injection

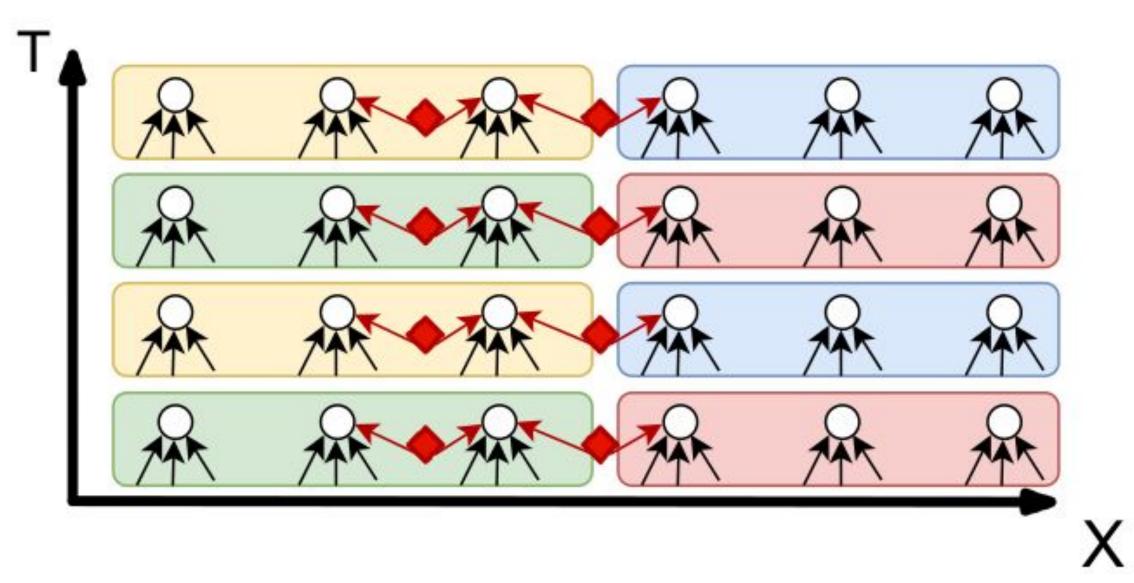
- Iterate over sources, each has 3-d coordinates
- Indirect accesses to scatter injection to neighbouring points
- Aligned in time, not in space

Listing 1: A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.

1 for t = 1 to nt do **for** x = 1 to nx do 2 **for** y = 1 to ny do 3 for z = 1 to nz do 4 $A(t, x, y, z) \equiv u[t, x, y, z] = u[t-1, x, y, z] + \sum_{r=1}^{r=so/2} w_r ($ 5 u[t-1, x - r, y, z] + u[t-1, x + r, y, z] + u[t-1, x, y - r, z] + u[t-1, x, y + r, z] + u[t-1, x, y, z - r] + u[t-1, x, y, z + r]);foreach s in sources do // For every source 6 for i = 1 to np do // Get the points affected 7 xs, ys, zs = map(s, i) / / through indirection8 u[t, xs, ys, zs] + = f(src(t, s)) / / add their impact9 on the field

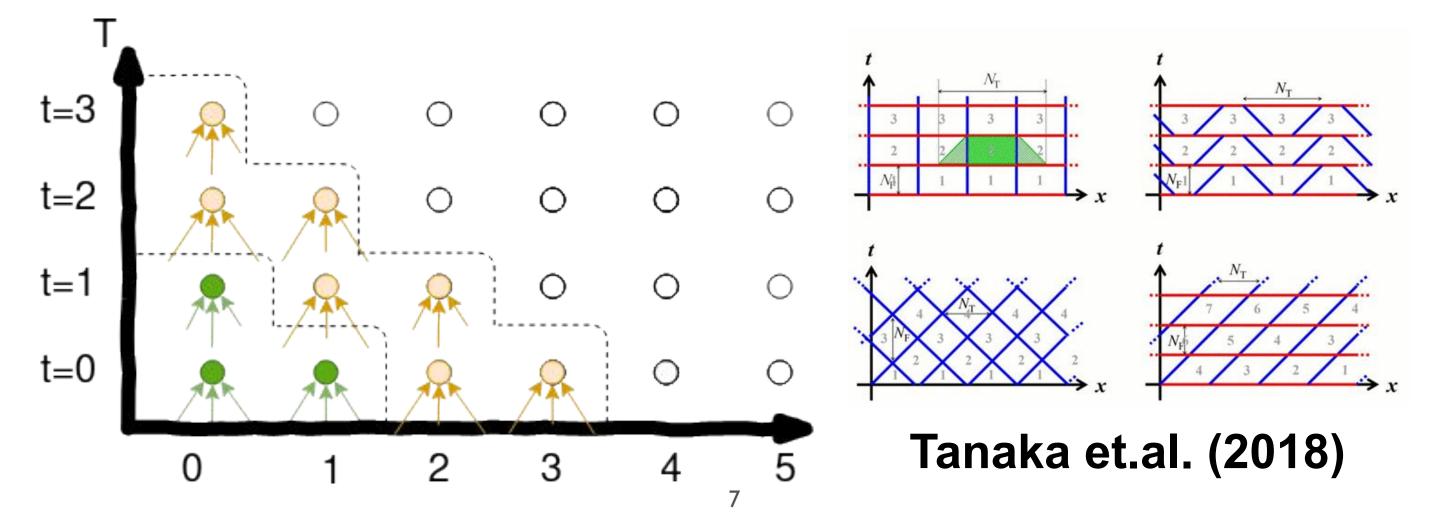
Applying space-blocking

- Spatial blocking:
 - Decompose grids into block tiles/ Partitioning iteration space to smaller chunks/blocks
 - Improve data locality => Increase performance (Rich literature)
 - Sparse off-the-grid operators are iterated as without blocking



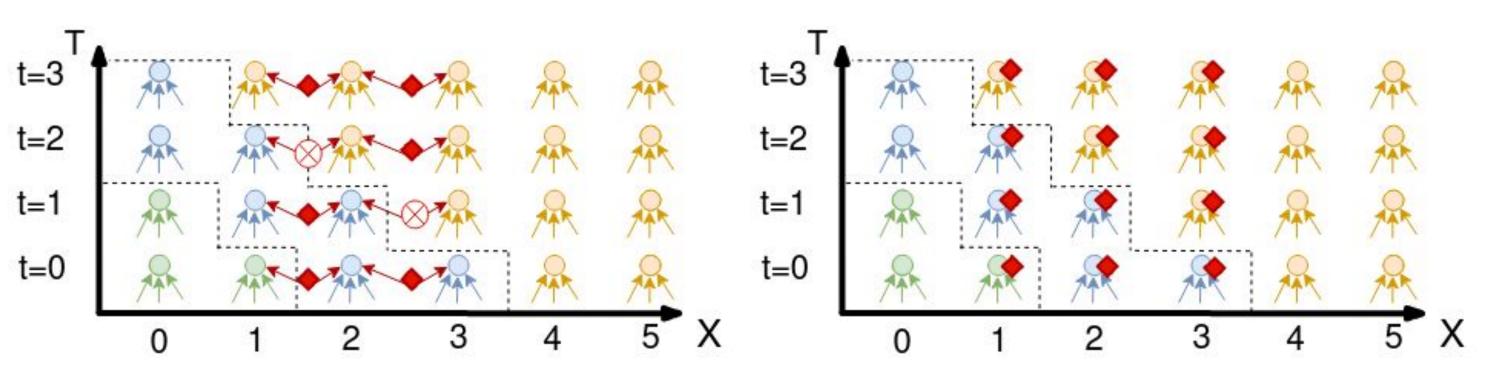
Applying temporal-blocking

- Temporal blocking:
 - Space blocking but extend reuse to time-dimension.
 - Update grid points in future where/when (space+time) possible
 - Rich literature, several variants of temporal blocking, shapes, schemes
 - -Wave-front / Skewing (Our POC approach)
 - -Diamonds, Trapezoids, Overlapped, Hybrid models



Off-the-grid operators: the issue

- Data dependences violations happen while a temporal update
- Source injection in a different iteration space
- When a sparse operator exists in the boundary between space-time blocks, order of updates is not preserved
- Solution: Need to align off-the-grid operators



Methodology

- A negligible-cost scheme to precompute the source injection contribution.
- Align source injection data dependences to the grid
- Applicable to other fields as well

Iterate over sources and store indices of affected points

- Inject to a **zero-initialized grid** for one (or a few more)
- **Hypothesis:** non-zero values at the first time-steps
- **Independent** of the injection and interpolation type (e.g. non-linear injection)

Listing 2: Source injection over an empty grid. No PDE stencil update is happening.

1	for $t = 1 to 2 do$
2	foreach s in sources do
3	for <i>i</i> = 1 <i>to</i> np do
	xs, ys, zs = map(s, i);
5	$\begin{vmatrix} \mathbf{for} & i = 1 \text{ to np do} \\ xs, ys, zs = map(s, i); \\ u[t, xs, ys, zs] + = f(src(t, s)) \end{vmatrix}$

• Then, we store the non-zero grid point coordinates

Generate sparse binary mask, unique IDs and decompose wavefields

Perform source injection to decompose the off-the-grid wavefields to on-the-grid per point wavefields.

	Off-the-grid	Aligned
len(sources)	n_src	n_aff_pts
len(sources.coords)	(n_src, 3)	(n_aff_pts, 3)
len(sources.data)	(n_src, nt)	(n_aff_pts, nt)

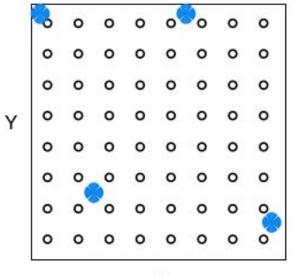
Listing 3: Decomposing the source injection wavefields.

1 for t = 1 to nt do

3

4 | |
$$xs, ys, zs = map(s, i);$$

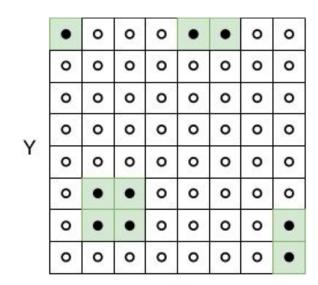
5
$$|$$
 src_dcmp[t, SID[xs, ys, zs]] + = $f(src(t, s);$



Х

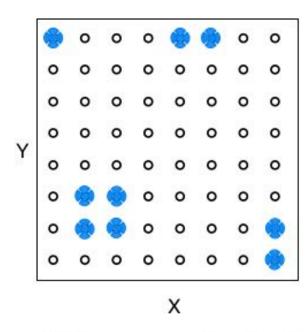
(a) Sources are sparsely distributed at off-the-grid positions.

1	0	0	0	2	3	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0
0	4	5	0	0	0	0	0
0	6	7	0	0	0	0	(
0	0	0	0	0	0	0	npts



Х

(b) Identify unique points affected (SM).



(c) Assign a unique ID to every affected point (SID).

(d) Sources are aligned with grid positions.

Y

Fuse iteration spaces

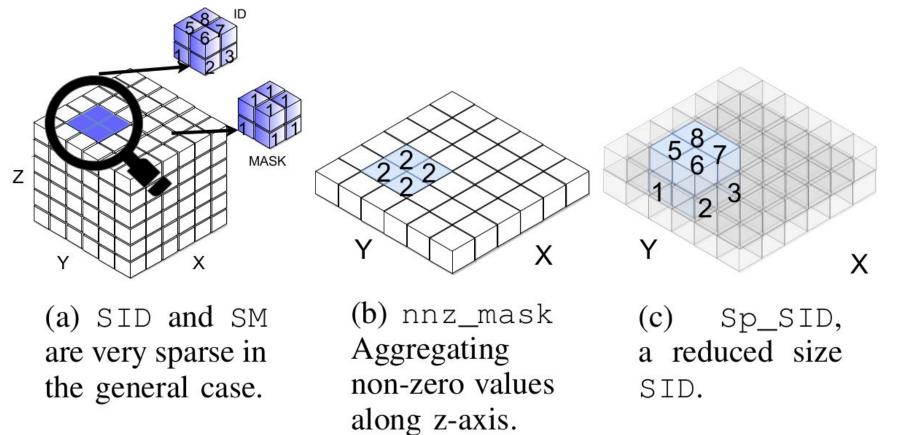
- Indirection mapping has changed. We still use indirections but now they are on the point.
- By using the aligned structure, we fuse the source injection loop inside the kernel update iteration space.
- The source mask SM is used to add (if 1) or not (if 0) the impact and SID is used to indirect to the impact values using the traversed grid coordinates.

Listing 4: Stencil kernel update with fused source injection.

1 for
$$t = 1$$
 to nt do
2 | for $x = 1$ to nx do
3 | for $y = 1$ to ny do
4 | for $z = 1$ to nz do
5 | | $A(t, x, y, z, s);$
6 | $| A(t, x, y, z, s);$
7 | $| u[t, x, y, z2] + = SM[x, y, z2] * src_dcmp[t, SID[x, y, z2]];$

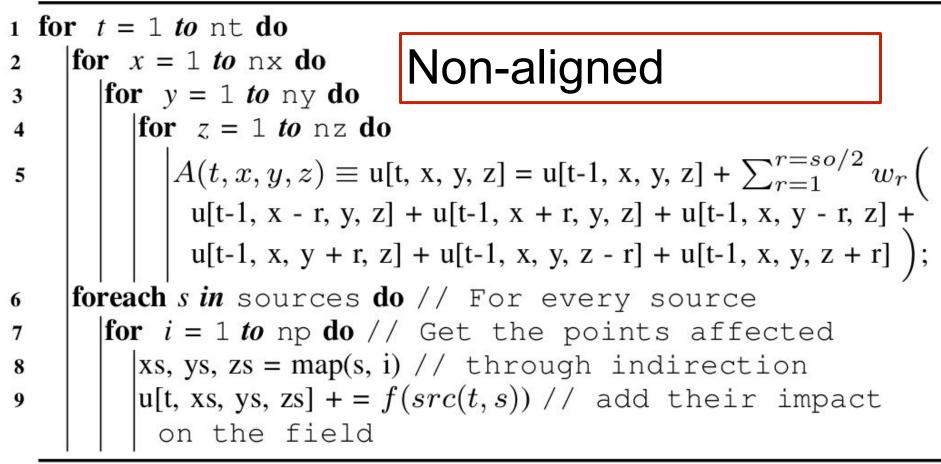
Reducing the iteration space size

- Perform only necessary operations
- Aggregate NZ along the z- axis keeping count of them in a structure named nnz_mask
- Reduce the size of SID by cutting off zero z-slices



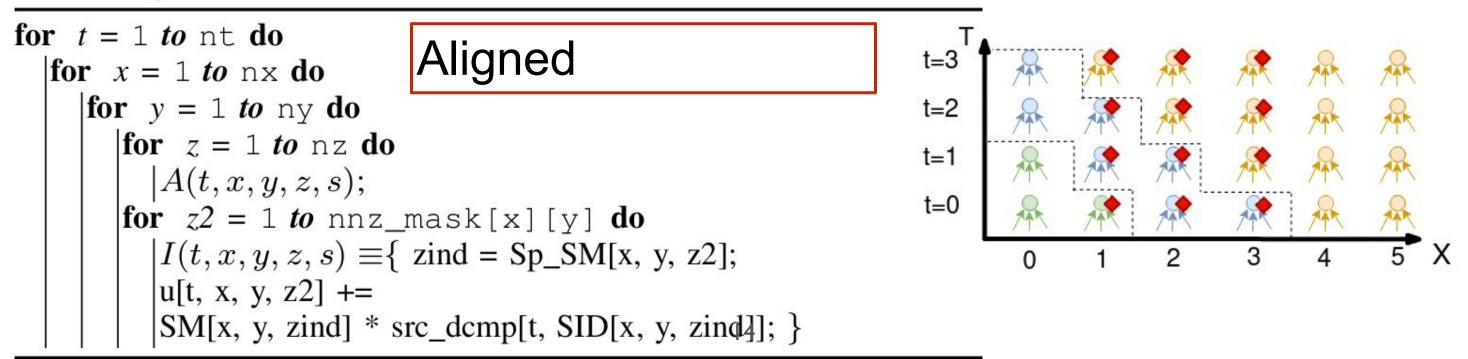
Listing 5: Stencil kernel update with reduced size iteration space for source injection.

1 for t = 1 to nt do 2 |for x = 1 to nx do 3 ||for y = 1 to ny do 4 ||for z = 1 to nz do 5 || ||A(t, x, y, z, s); 6 || ||A(t, x, y, z, s); 6 || || I(t, x, y, z, s) \equiv { zind = Sp_SID[x, y, z2]; 8 || || || I(t, x, y, z2] += src_dcmp[t, SID[x, y, zind]]; } **Listing 1:** A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.



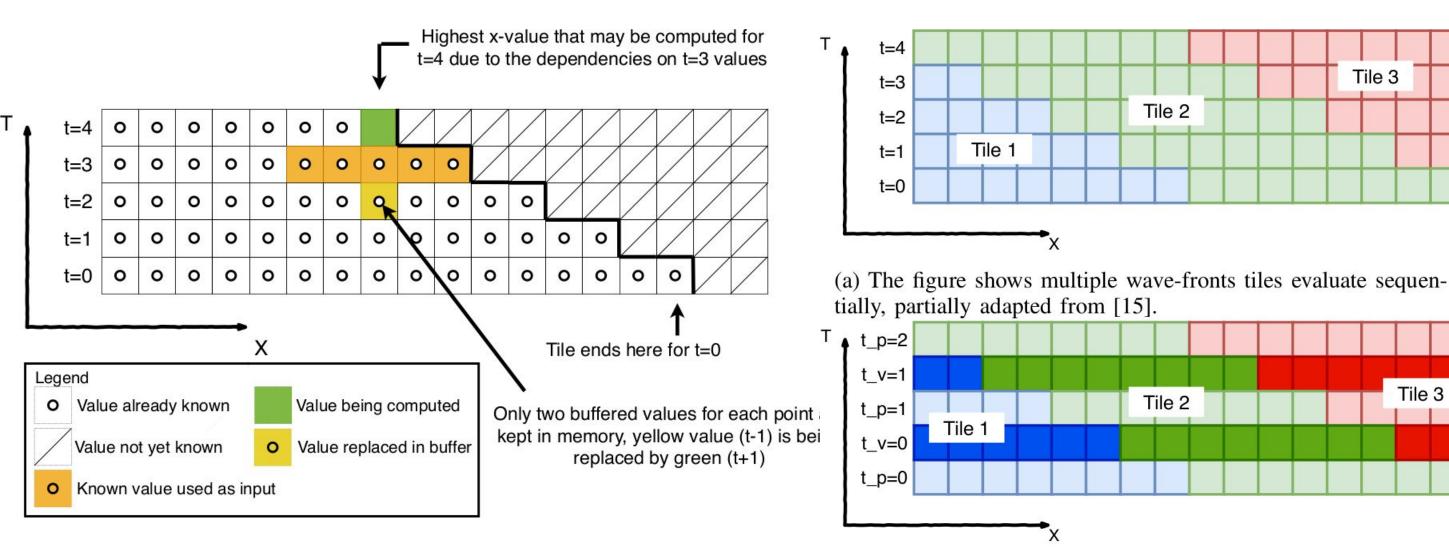
Aligned to grid
Same OPS
Parallelism
SIMD
Apply TB

Listing 5: Stencil kernel update with fused - reduced size iteration space - source injection.



Applying wave-front temporal blocking

- Aligning, automated in Devito DSL; TB with manual loop transformation
- The new source coordinates are aligned to the grid now
- Skewing factor depends on data dependency distances



YASK, Yount et. al (2016)

(b) The figure shows multiple wave-front tiles evaluated sequentially in multigrid stencil codes.

Tile 3

Listing 6: The figure shows the loop structure after applying our proposed scheme.

```
for t_tile in time_tiles do
  for xtile in xtiles do
    for ytile in ytiles do
      for t in t_tile do
         OpenMP parallelism
         for xblk in xtile do
           for yblk in ytile do
              for x in xblk do
                for y in yblk do
                  SIMD vectorization
                  for z = 1 to nz do
                    |A(t, x - time, y - time, z, s);
                  for z2 = 1 to nnz_mask[x][y] do
                    |I(t, x - time, y - time, z2, s);
```

Experimental evaluation: the models

Isotropic Acoustic

Generally known, single scalar PDE, laplacian like, low cost

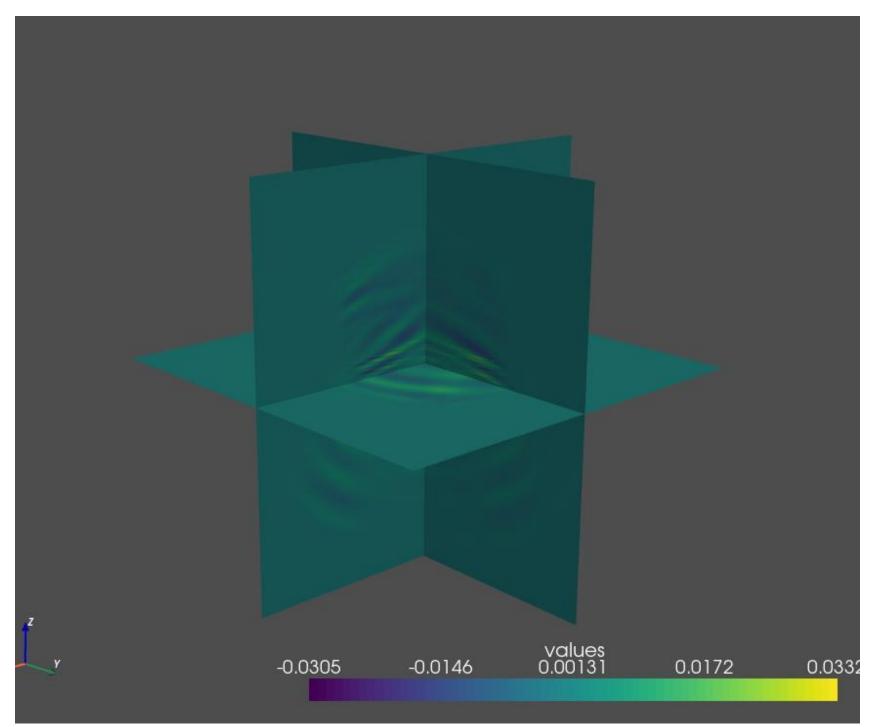
• Isotropic Elastic

Coupled system of a vectorial and tensorial PDE, explosive source, increased data movement, first order in time, cross-loop data dependences

Anisotropic Acoustic

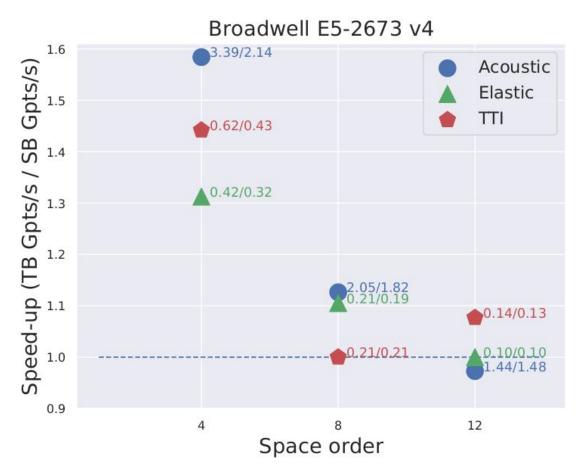
Industrial applications, rotated laplacian, coupled system of two scalar PDEs

Industrial-level, 512^3 grid points, 512ms simulation time, damping fields ABCs



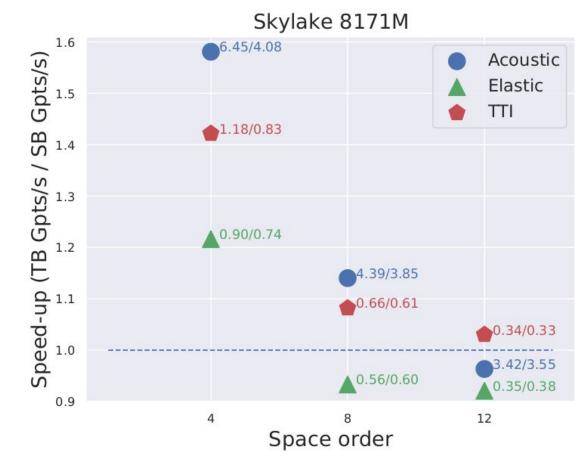
Velocity field, TTI wave propagation after 512ms

Experimental evaluation: the results



(a) Throughput speed-up of kernels for Broadwell.

Azure model Architecture	E16s v3 Broadwell	E32s v3 Skylake
vCPUs	16	32
GiB memory	128	256
Model name	E5-2673 v4	8171M
CPUs	16	32
Thread(s) per core	2	2
Core(s) per socket	8	16
Socket(s)	1	1
NUMA node(s)	1	1
Model	79	85
CPU MHz	2300	2100
L1d cache	32K	32K
L1i cache	32K	32K
L2 cache	256K	1024K
L3 cache	51200K	36608K

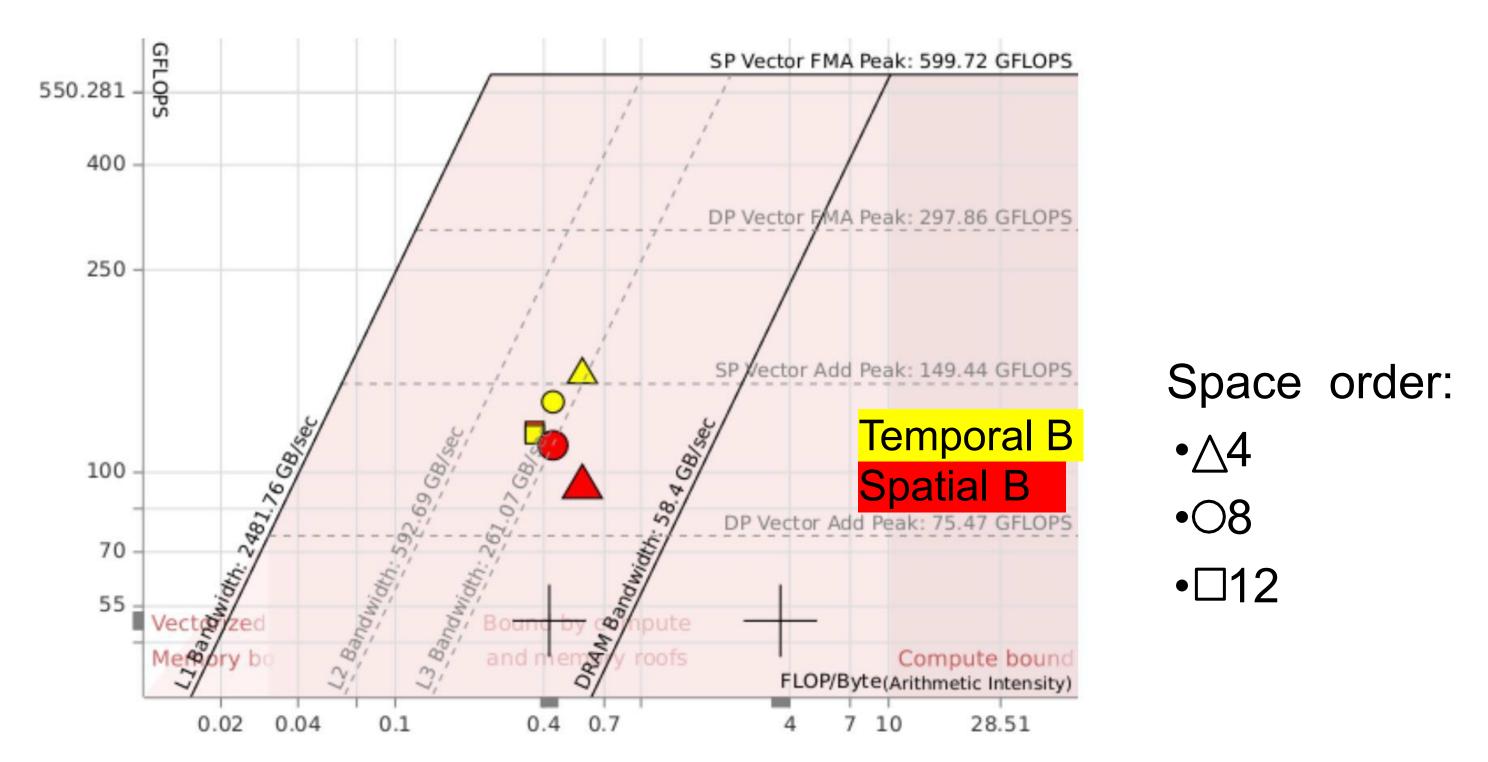


(b) Throughput speed-up of kernels for Skylake.

- Benchmark on Azure VMs
- GCC, ICC
- Thread pinning
- OpenMP, SIMD
- Aggressive auto-tuning

TABLE I: VM specification

Cache-aware roofline model



Broadwell, isotropic acoustic, 512^3 grid points, 512ms

Conclusions

- We presented an approach to apply temporal blocking on stencil kernels with sparse off-the-grid operators.
- The additional cost is negligible compared to the achieved gains.
- Solution built on top of Devito-DSL
- Performance gains of up to 1.6x on low order (4) and 1.2x on medium order (8).

Work presented is inherited from: Bisbas, G., Luporini, F., Louboutin, M., Nelson, R., Gorman, G., & Kelly, P.H. (2020). Temporal blocking of finite-difference stencil operators with sparse "off-the-grid" sources. Accepted at IPDPS'21. Available online: https://arxiv.org/abs/2010.10248



• Open-source, on top of Devito v4.2.3 https://github.com/georgebisbas/devito

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Engineering and Physical Sciences

- Integration/ Automation
- GPUs
- High-order stencils

Website: http://www.devitoproject.org GitHub: https://github.com/devitocodes/devito Slack: https://opesci-slackin.now.sh



Acknowledgements

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- Navjot Kukreja (Imperial College)
- John Washbourne (Chevron)
- Edward Caunt (Imperial College)

Thank you for your attention! Questions?

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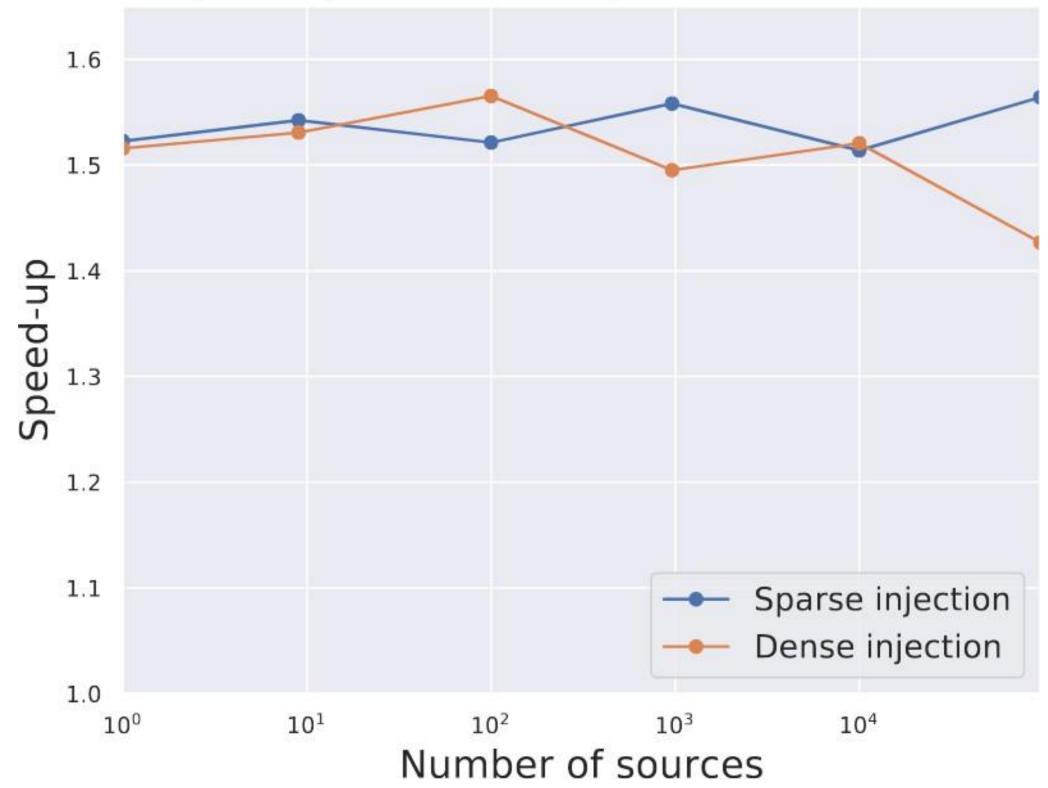


References

- Bisbas, G., Luporini, F., Louboutin, M., Nelson, R., Gorman, G., & Kelly, P.H. (2020). Temporal blocking of finite-difference stencil operators with sparse "off-the-grid" sources. Accepted at IPDPS'21. Available online: https://arxiv.org/abs/2010.10248
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- Louboutin, M., M., Lange, F., Luporini, N., Kukreja, P. A., Witte, F. J., Herrmann, P., Velesko, and G. J., Gorman. "Devito (v3.1.0): an embedded domain-specific language for finite differences and geophysical exploration". Geoscientific Model Development 12, no.3 (2019): 1165–1187.
- Yount, C., & Duran, A. (2016). Effective Use of Large High-Bandwidth Memory Caches in HPC Stencil Computation via Temporal Wave-Front Tiling. (2016) 7th International Workshop on Performance Modeling, Benchmarking and Simulation of High Performance Computer Systems (PMBS), 65-75.

Corner cases, increasing number of sources

Speed-up for increasing number of sources



The generated C code - stencil update

```
#pragma omp tor collapse(1) schedule(dynamic,1)
             for (int x0_blk0 = x_m; x0_blk0 <= x_M; x0 blk0 += x0 blk0 size)</pre>
             ſ
                    for (int y0 blk0 = y m; y0 blk0 <= y M; y0 blk0 += y0 blk0 size)</pre>
                          for (int x = x0 blk0; x <= x0 blk0 + x0 blk0 size - 1; x += 1)</pre>
                                 for (int y = y0 blk0; y <= y0 blk0 + y0 blk0 size - 1; y += 1)</pre>
                                       #pragma omp simd aligned(damp,uref,vp:32)
                                       for (int z = z_m; z <= z_M; z += 1)</pre>
                                       {
                                              float r14 = -2.84722222F*uref[t1][x + 8][y + 8][z + 8];
                                              float r13 = 1.0/dt;
                                              float r12 = 1.0/(dt*dt);
                                             float r11 = 1.0/(vp[x + 8][y + 8][z + 8]*vp[x + 8][y + 8][z + 8]);
                                             uref[t0][x + 8][y + 8][z + 8] = (r11*(-r12*(-2.0F*uref[t1][x + 8][y + 8][z + 8] + 8][z + 8][z + 8]] + 8[z + 8][z + 8][z
 uref[t_2][x + 8][y + 8][z + 8]) + r_{13}(damp[x + 1][y + 1][z + 1]*uref[t_1][x + 8][y + 8][z + 8])
 (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 4] + uref[t1][x + 8][y + 8][z + 12]) +
 2.53968254e-2F*(uref[t1][x + 8][y + 8][z + 5] + uref[t1][x + 8][y + 8][z + 11]) -
 2.0e-1F*(uref[t1][x + 8][y + 8][z + 6] + uref[t1][x + 8][y + 8][z + 10]) + 1.6F*(uref[t1][x + 8]
 [y + 8][z + 7] + uref[t1][x + 8][y + 8][z + 9]))/((h_z*h_z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h_z*h_z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h_z*h_z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][z + 9]))/((h_z*h_z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][y + 8][z + 9]))/((h_z*h_z)) + (r14 - 1.78571429e-3F*(uref[t1][x + 8][y + 8][
 + 8][y + 4][z + 8] + uref[t1][x + 8][y + 12][z + 8]) + 2.53968254e-2F*(uref[t1][x + 8][y + 5][z])
 B] + uref[t1][x + 8][y + 11][z + 8]) - 2.0e-1F*(uref[t1][x + 8][y + 6][z + 8] + uref[t1][x + 8][
 + 10[z + 8] + 1.6F*(uref[t1][x + 8][y + 7][z + 8] + uref[t1][x + 8][y + 9][z + 8]))/((h y*h y))
 + (r14 - 1.78571429e-3F*(uref[t1][x + 4][y + 8][z + 8] + uref[t1][x + 12][y + 8][z + 8]) +
 2.53968254e-2F*(uref[t1][x + 5][y + 8][z + 8] + uref[t1][x + 11][y + 8][z + 8]) -
 2.0e-1F*(uref[t1][x + 6][y + 8][z + 8] + uref[t1][x + 10][y + 8][z + 8]) + 1.6F*(uref[t1][x + 7]
 [y + 8][z + 8] + uref[t1][x + 9][y + 8][z + 8]))/((h x*h x)))/(r11*r12 + r13*damp[x + 1][y + 1][
 + 1]);
}
}

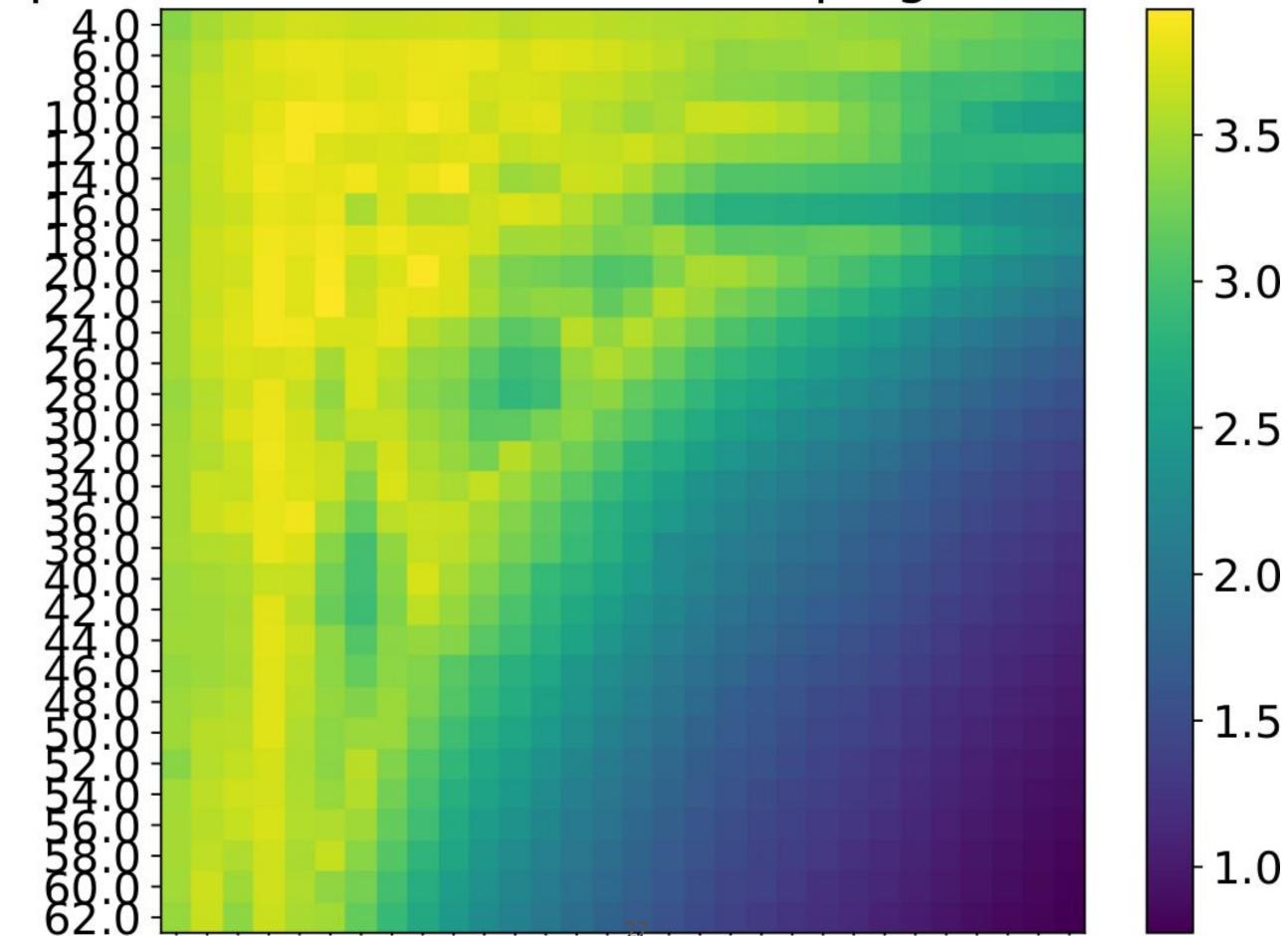
                                     }
```

The generated C code - source injection

```
/* Begin section1 */
   #pragma omp parallel num threads(nthreads nonaffine)
     int chunk size = (int)(fmax(1, (1.0F/3.0F)*(p src M - p src m + 1)/nthreads nonaffine));
     #pragma omp for collapse(1) schedule(dynamic,chunk size)
     for (int p src = p src m; p src <= p src M; p src += 1)</pre>
       int ii src 0 = (int)(floor((-o x + src coords[p src][0])/h x));
       int ii src 1 = (int)(floor((-o y + src coords[p src][1])/h y));
       int ii src 2 = (int)(floor((-o z + src coords[p src][2])/h z));
       int ii src 3 = (int)(floor((-o z + src coords[p src][2])/h z)) + 1;
       int ii src 4 = (int)(floor((-o y + src coords[p_src][1])/h_y)) + 1;
       int ii src 5 = (int)(floor((-o x + src coords[p src][0])/h x)) + 1;
       float px = (float)(-h x*(int)(floor((-o x + src coords[p src][0])/h x)) - o x + src coords[p src][0]);
       float py = (float)(-h_y*(int)(floor((-o_y + src_coords[p_src][1])/h_y)) - o_y + src_coords[p_src][1]);
       float pz = (float)(-h z*(int)(floor((-o z + src coords[p src][2])/h z)) - o z + src coords[p src][2]);
       if (ii src 0 >= x m - 1 && ii src 1 >= y m - 1 && ii src 2 >= z m - 1 && ii src 0 <= x M + 1 && ii src 1
<= y M + 1 && ii src 2 <= z M + 1)
         float r0 = 4.49016082216644F*(vp[ii_src_0 + 8][ii_src_1 + 8][ii_src_2 + 8]*vp[ii_src_0 + 8][ii_src_1 + 8]
[ii_src_2 + 8] (-px*py*pz/(h_x*h_y*h_z) + px*py/(h_x*h_y) + px*pz/(h_x*h_z) - px/h_x + py*pz/(h_y*h_z) - py/h_y -
pz/h z + 1)*src[time][p src];
         #pragma omp atomic update
         uref[t0][ii src 0 + 8][ii src 1 + 8][ii src 2 + 8] += r0;
       }
       if (ii src 0 >= x m - 1 && ii src 1 >= y m - 1 && ii src 3 >= z m - 1 && ii src 0 <= x M + 1 && ii src 1
<= y M + 1 && ii src 3 <= z M + 1)
       {
         float r1 = 4.49016082216644F*(vp[ii src 0 + 8][ii src 1 + 8][ii src 3 + 8]*vp[ii src 0 + 8][ii src 1 + 8]
[ii_src_3 + 8])*(px*py*pz/(h_x*h_y*h_z) - px*pz/(h_x*h_z) - py*pz/(h_y*h_z) + pz/h_z)*src[time][p_src];
         #pragma omp atomic update
         uref[t0][ii src 0 + 8][ii src 1 + 8][ii src 3 + 8] += r1;
       }
       if (ii src 0 >= x m - 1 && ii src 2 >= z m - 1 && ii src 4 >= y m - 1 && ii src 0 <= x M + 1 && ii src 2
<= z M + 1 && ii src 4 <= y M + 1)
       ł
         float r2 = 4.49016082216644F*(vp[ii src 0 + 8][ii src 4 + 8][ii src 2 + 8]*vp[ii src 0 + 8][ii src 4 + 8]
[ii src 2 + 8])*(px*py*pz/(h x*h y*h z) - px*py/(h x*h y) - py*pz/(h y*h z) + py/h y)*src[time][p src]:
```

		Eloor, ceil of
	Algorithm 3: Source injection pseudocode.	·
1	for t = 1 to nt do	off- the-grid
2	foreach s in sources do	
3	# Find on the grid coordinates	
4	$src_x_min = floor(src_coords[s][0], ox)$	
5	$src_x_max = ceil(src_coords[s][0], ox)$	
6	$src_y_min = floor(src_coords[s][1], oy)$	
7	$src_y_max = ceil(src_coords[s][1], oy)$	Waights of
8	$src_z_min = floor(src_coords[s][2], oz)$	Weights of
9	$src_z_max = ceil(src_coords[s][2], oz)$	impact
10	# Compute weights	
11	$px = f(src_coords[s][0], ox)$	
12	$py = f(src_coords[s][1], oy)$	
13	$pz = f(src_coords[s][2], oz)$	Unrolled loop
14	# Unrolled for 8 points	rid then forint ach
15	if src_x_min, src_y_min, src_z_min in gr	
16	$r0 = v(src_x_min, src_y_min, src_z_min)$	in, src[t][s])
17	u[t, src_x_min, src_y_min, src_z_min]	$+ = r^{(0)}$ and dd to field
	:	a
18	if src_x_max, src_y_max, src_z_max in g	
19	$r7 = v(src_x_max, src_y_max, src_z_max)$	
20	u[t, src_x_max, src_y_max, src_z_max]	J + = r/)

Gpts/s for fixed tile size. (Sweeping block sizes)



Algorithm 3: Source injection pseudocode.

```
1 for t = l to nt do
     foreach s in sources do
2
        # Find on the grid coordinates
3
        src_x_min = floor(src_coords[s][0], ox)
4
        src_x_max = ceil(src_coords[s][0], ox)
5
        # Compute weights
6
7
        px = f(src\_coords[s][0], ox)
        # Unrolled for 8 points (2^3, 3D \text{ case})
8
        if src_x_min, ... in grid then
9
           r0 = v(src_x_min, \dots src[t][s]);
10
           u[t, src_x_min, ...] + = r0)
11
        if src_x_max, ... in grid then
12
           r7 = v(src_x_max, \dots src[t][s]);
13
           u[t, src_x_max, ...] + = r7)_{28}
14
```

Cache aware roofline model

From here: https://crd.lbl.gov/departments/computer-science/par/research/roofline/introduction/

Effects of Cache Behavior on Arithmetic Intensity

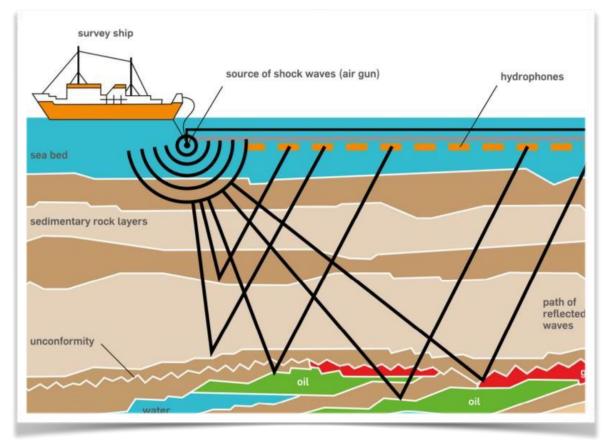
The Roofline model requires an estimate of total data movement. On cache-based architectures, the 3C's cache model highlights the fact that there can be more than simply compulsory data movement. Cache capacity and conflict misses can increase data movement and reduce arithmetic intensity.

Similarly, superfluous cache write-allocations can result in a doubling of data movement. The vector initialization operation x[i]=0.0 demands one write allocate and one write back per cache line touched. The write allocate is superfluous as all elements of that cache line are to be overwritten. Unfortunately, the presence of hardware stream prefetchers can make it very difficult to quantify how much beyond compulsory data movement actually occurred.

A bit of background

- **PDEs** are everywhere: computational fluid dynamics, image processing, weather forecasting, seismic and medical imaging.
- Numerical analysis => finite-difference (FD) methods to solve DEs by approximating derivatives with finite differences.
- Devito: Fast Stencil Computation from Symbolic Specification
- Goal:

To improve performance of stencils stemming from practical applications using temporal blocking



Algorithm 3: Source injection pseudocode.		
1	or $t = 1 to$ nt do	
	foreach s in sources do	

Find on the grid coordinates

1

2

3

4

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src_x_min = floor(src_coords[s][0], ox)

src_x_max = ceil(src_coords[s][0], ox)

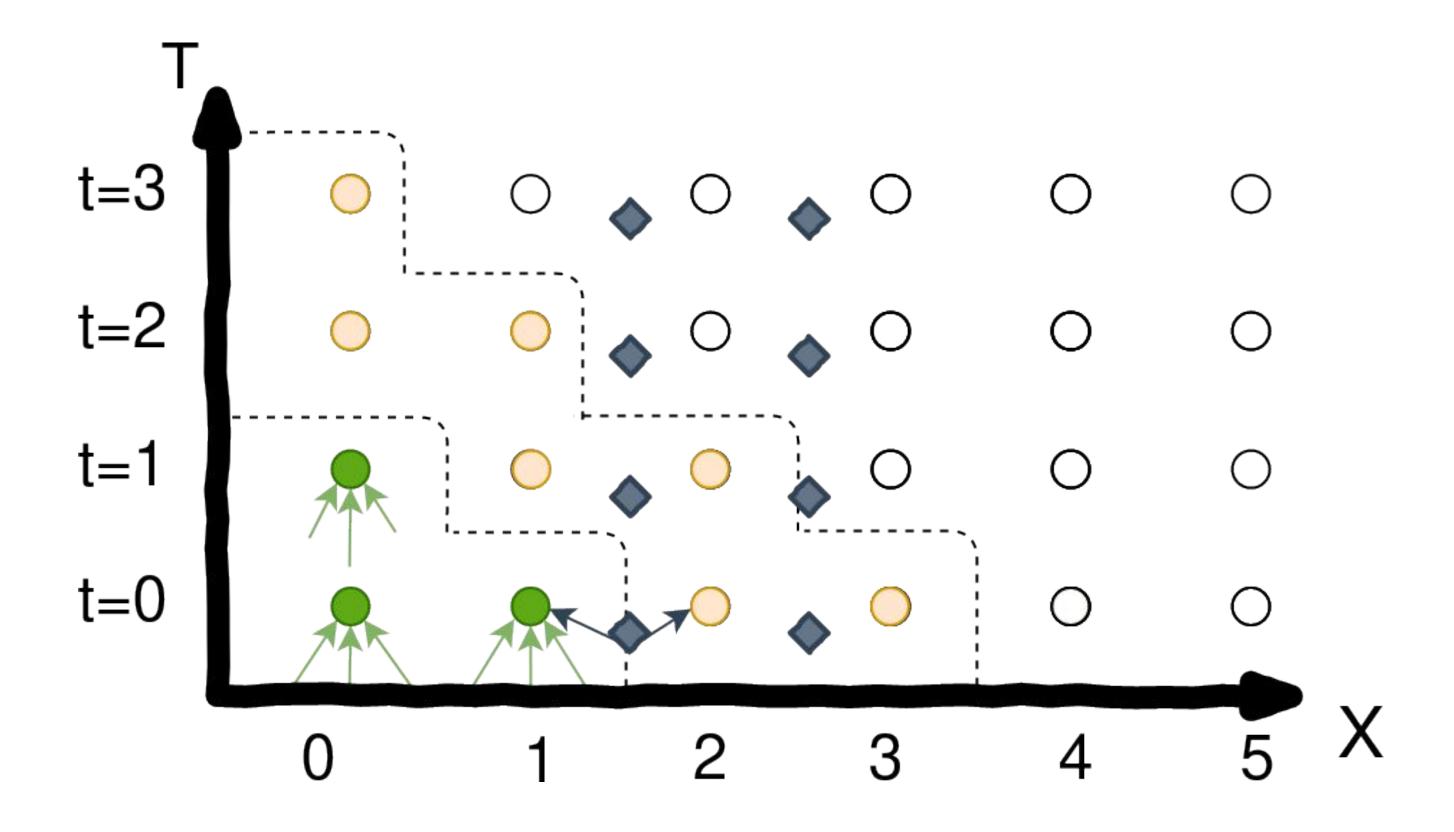
Compute weights
px = f(src_coords[s][0], ox)

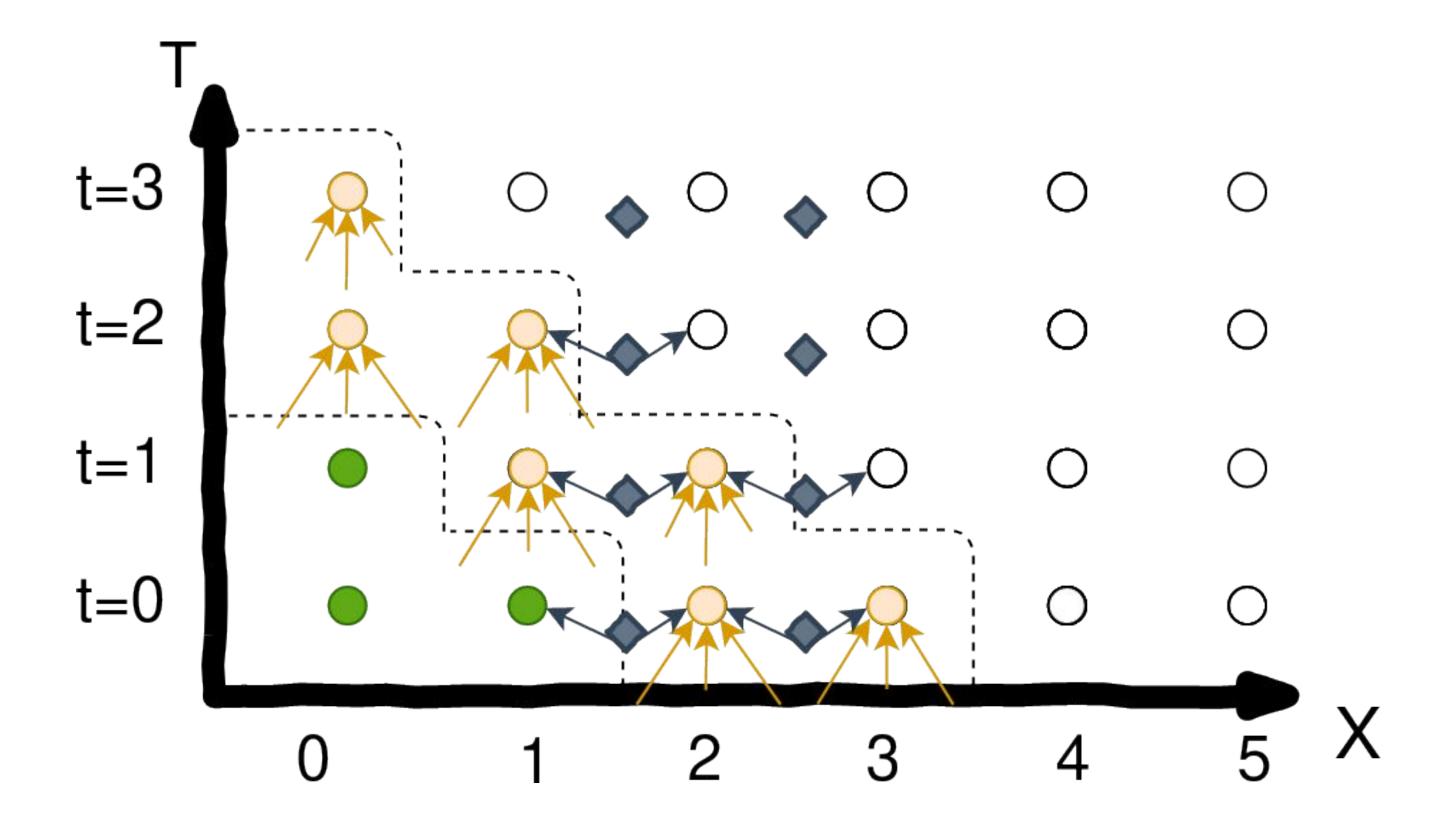
Unrolled for 8 points (2³, 3D case) if src_x_min , ... in grid then $|r0 = v(src_x_min, ... src[t][s]);$ $u[t, src_x_min, ...] + = r0)$

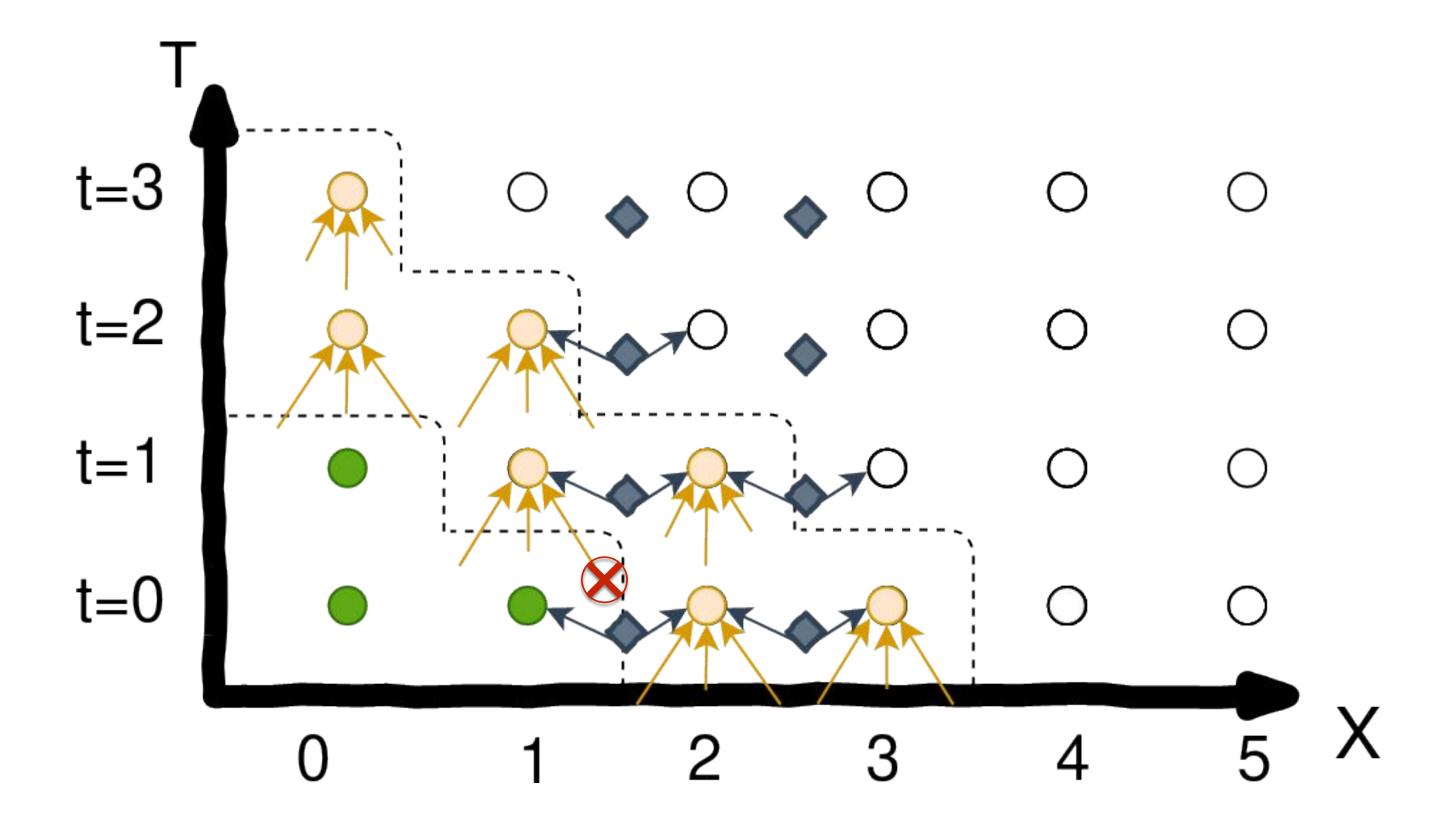
if src_x_max , ... in grid then $|r7 = v(src_x_max, ... src[t][s]);$ $u[t, src_x_max, ...] + = r7)$ Discover affected points

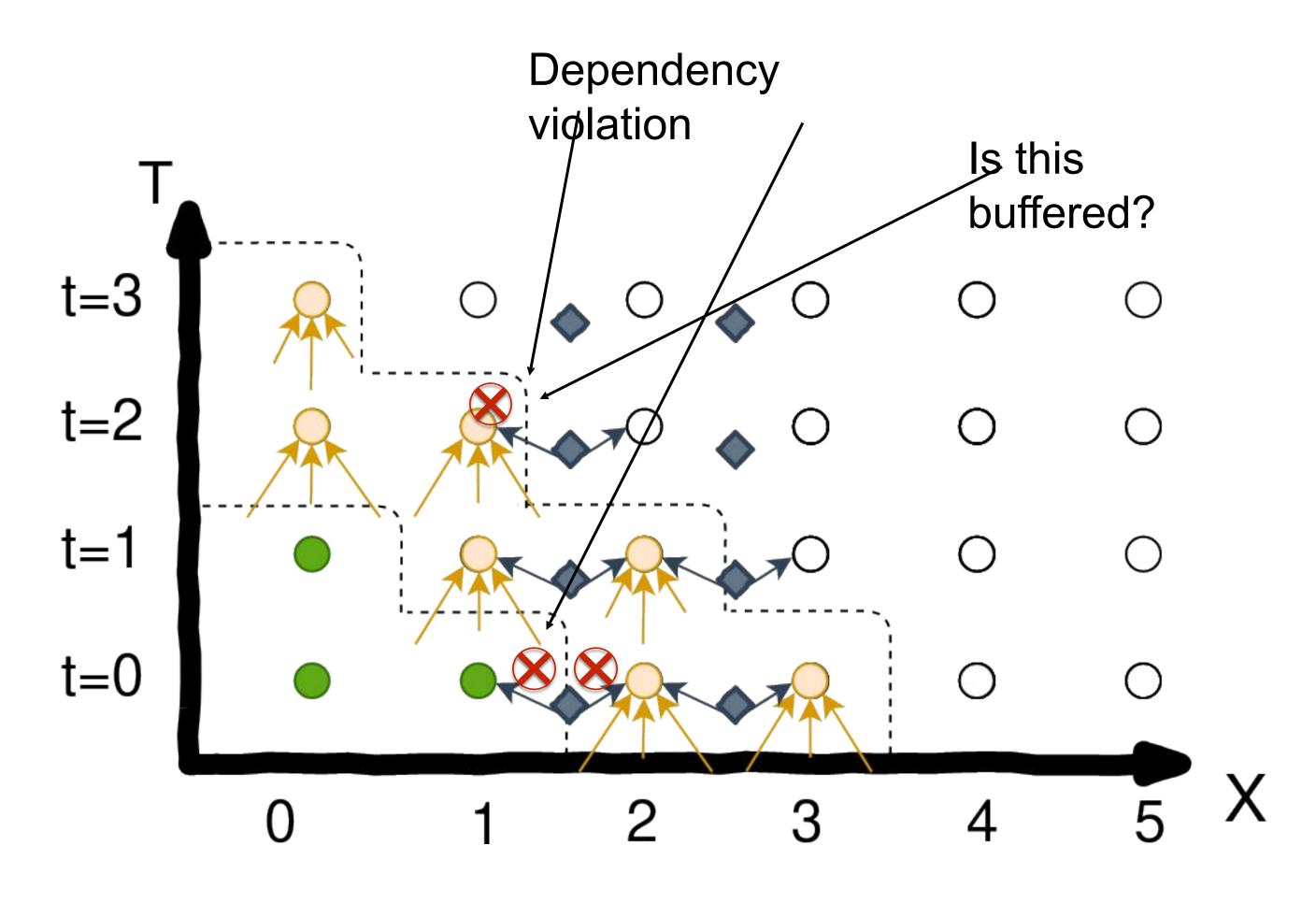
Weights of impact

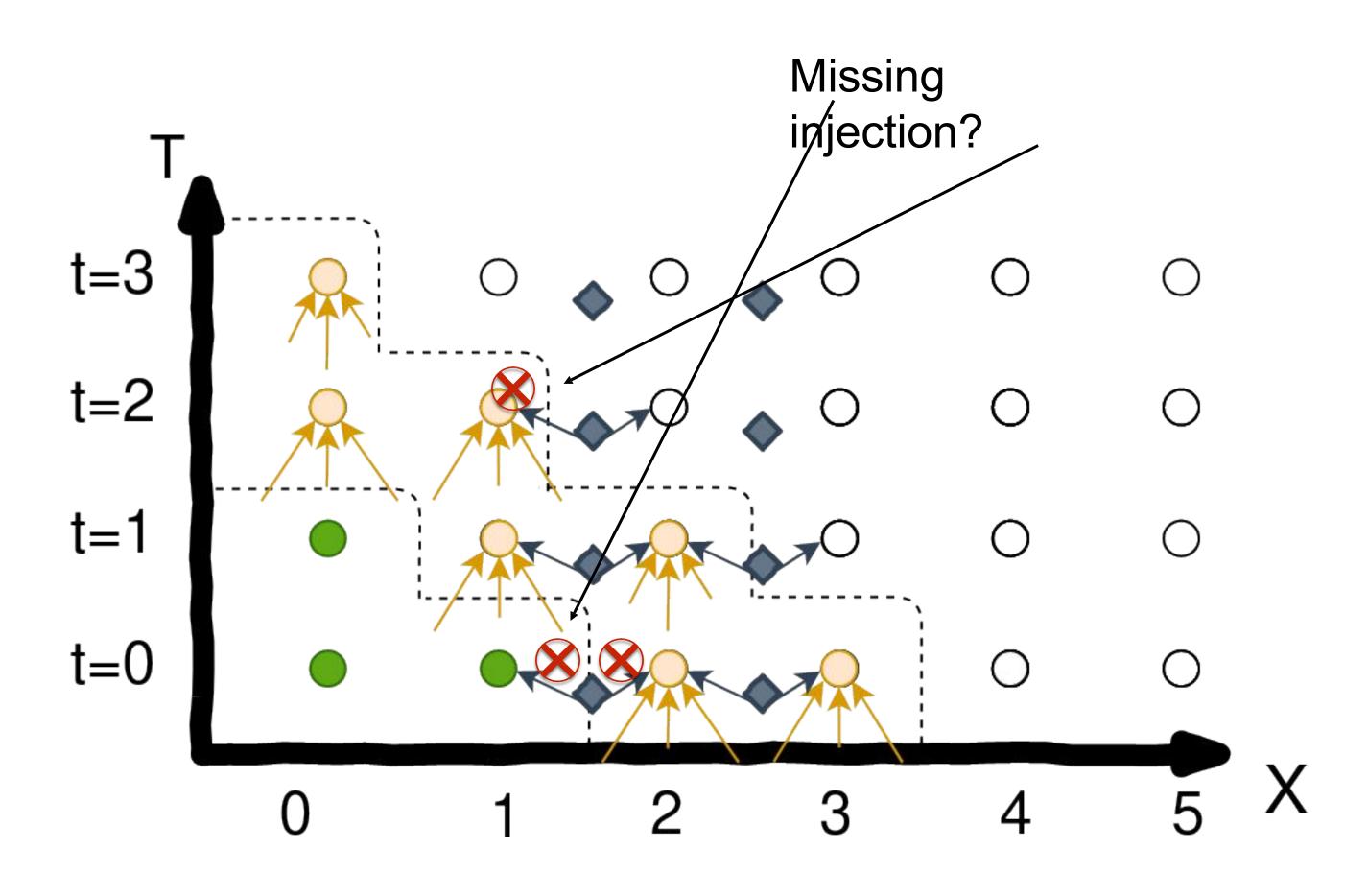
Unrolled loop for each affected point, compute injection part and add to field

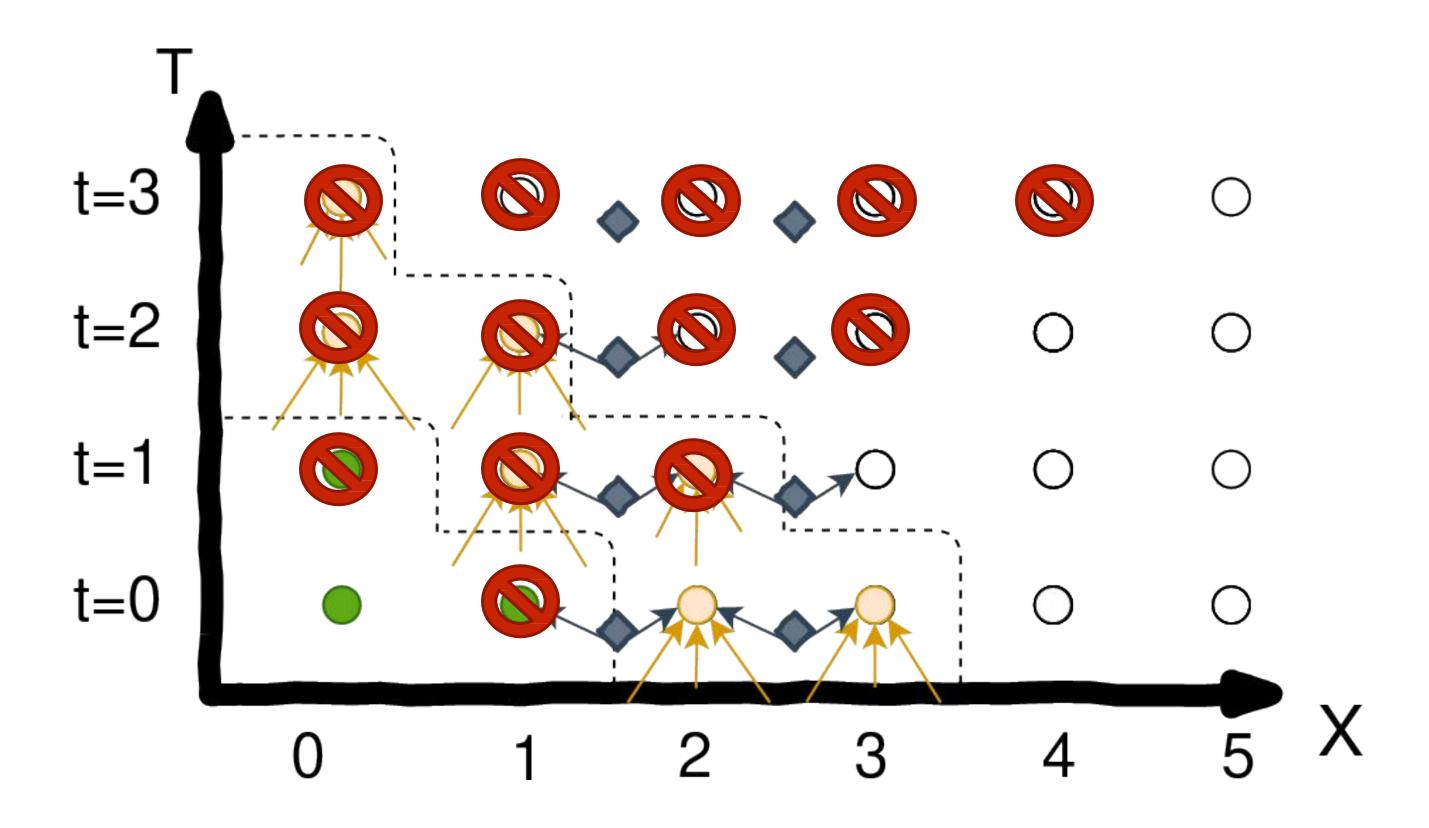












Algorithm 1: A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.

1 for t = 1 to nt do for x = 1 to nx do 2 for y = 1 to ny do 3 for z = 1 to nz do 4 $\begin{vmatrix} A(t, x, y, z) \equiv u[t, x, y, z] = u[t-1, x, y, z] + \sum_{r=1}^{r=so/2} w_r \\ u[t-1, x - r, y, z] + u[t-1, x + r, y, z] + u[t-1, x, y - r, z] + \end{vmatrix}$ 5 u[t-1, x, y + r, z] + u[t-1, x, y, z - r] + u[t-1, x, y, z + r]; foreach s in sources do 6 for *i* = 1 *to* np do 7 xs, ys, zs = map(s, i); 8 u[t, xs, ys, zs] + = f(src(t, s))9

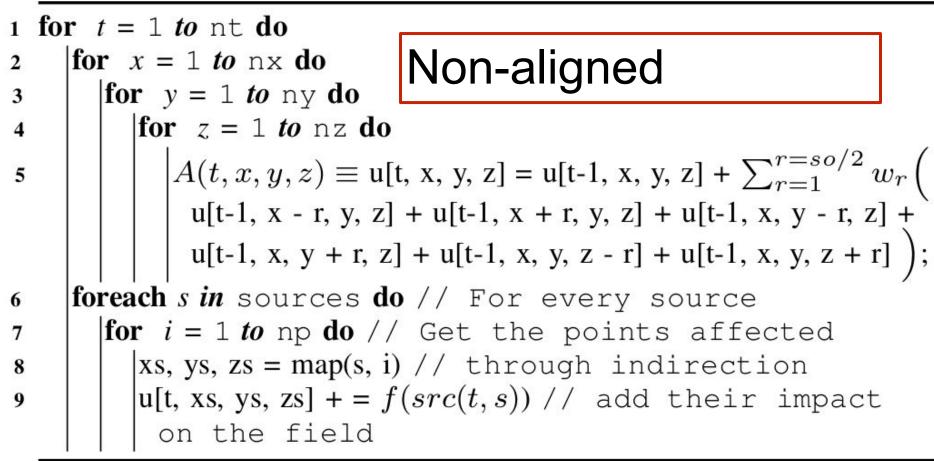
Non-aligned

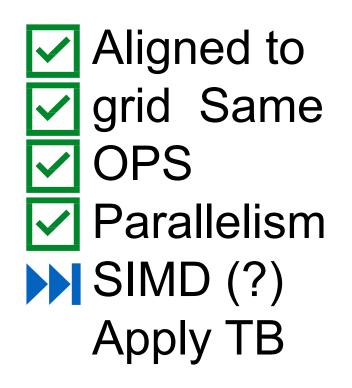
Algorithm 6: Stencil kernel update with fused - reduced size iteration space - source injection.

```
for t = 1 to nt do
  for x = 1 to nx do
    for y = 1 to ny do
       for z = 1 to nz do
         |A(t, x, y, z, s);
       for z^2 = 1 to nnz_mask[x][y] do
          zind = Sp_SM[x, y, z];
          u[t, x, y, z2] +=
          SM[x, y, zind] * src_dcmp[t, SID[x, y, zind]];
```

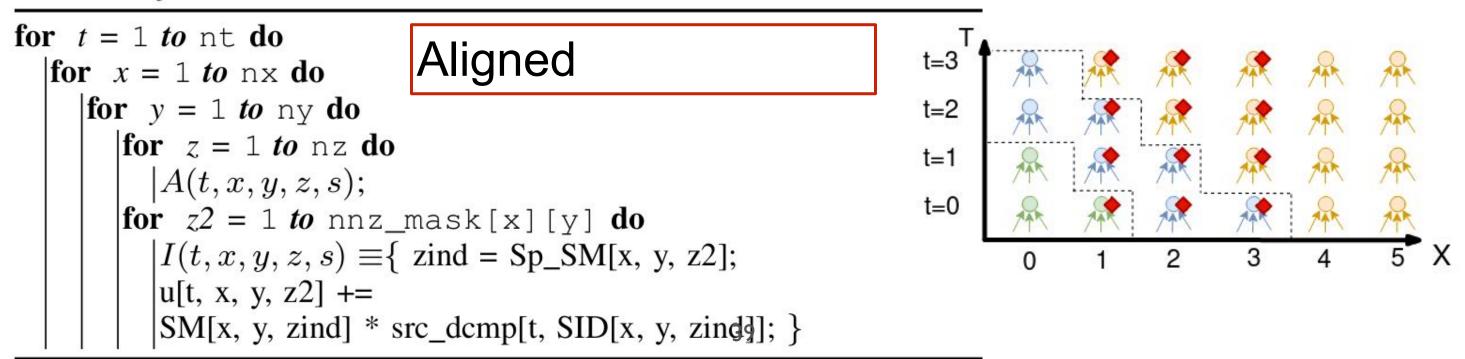
Aligne Ω

Listing 1: A typical time-stepping loop nest structure for a stencil update with source injection. This stencil has one temporal and three spatial dimensions.





Listing 5: Stencil kernel update with fused - reduced size iteration space - source injection.



The transformation in Devito-DSL

u = TimeFunction(name="u", grid=model.grid, space_order=so, time_order=2)
src_term = src.inject(field=u.forward, expr=src * dt**2 / model.m)
pde = model.m * u.dt2 - u.laplace + model.damp * u.dt
stencil = Eq(u.forward, solve(pde, u.forward))
op = Operator([stencil, src_term])

The transformation in

f : perform Sevite On Sempty grid

f = TimeFunction(name="f", grid=model.grid, space_order=so, time_order=2)
src_f = src.inject(field=f.forward, expr=src * dt**2 / model.m)
op_f = Operator([src_f])
op_f_sum = op_f.apply(time=3)

nzinds = np.nonzero(f.data[O]) # nzinds is a tuple

eq0 = Eq(sp_zi.symbolic_max, nnz_sp_source_mask[x, y] - 1, implicit_dims=(time, x, y)) eq1 = Eq(zind, sp_source_mask[x, y, sp_zi], implicit_dims=(time, x, y, sp_zi))

```
mask_expr = source_mask[x, y, zind] * save_src[time, source_id[x, y, zind]]
eq2 = Inc(usol.forward[t+1, x, y, zind], mask_expr, implicit_dims=(time, x, y, sp_zi))
```

pde_2 = model.m * usol.dt2 – usol.laplace + model.damp * usol.dt

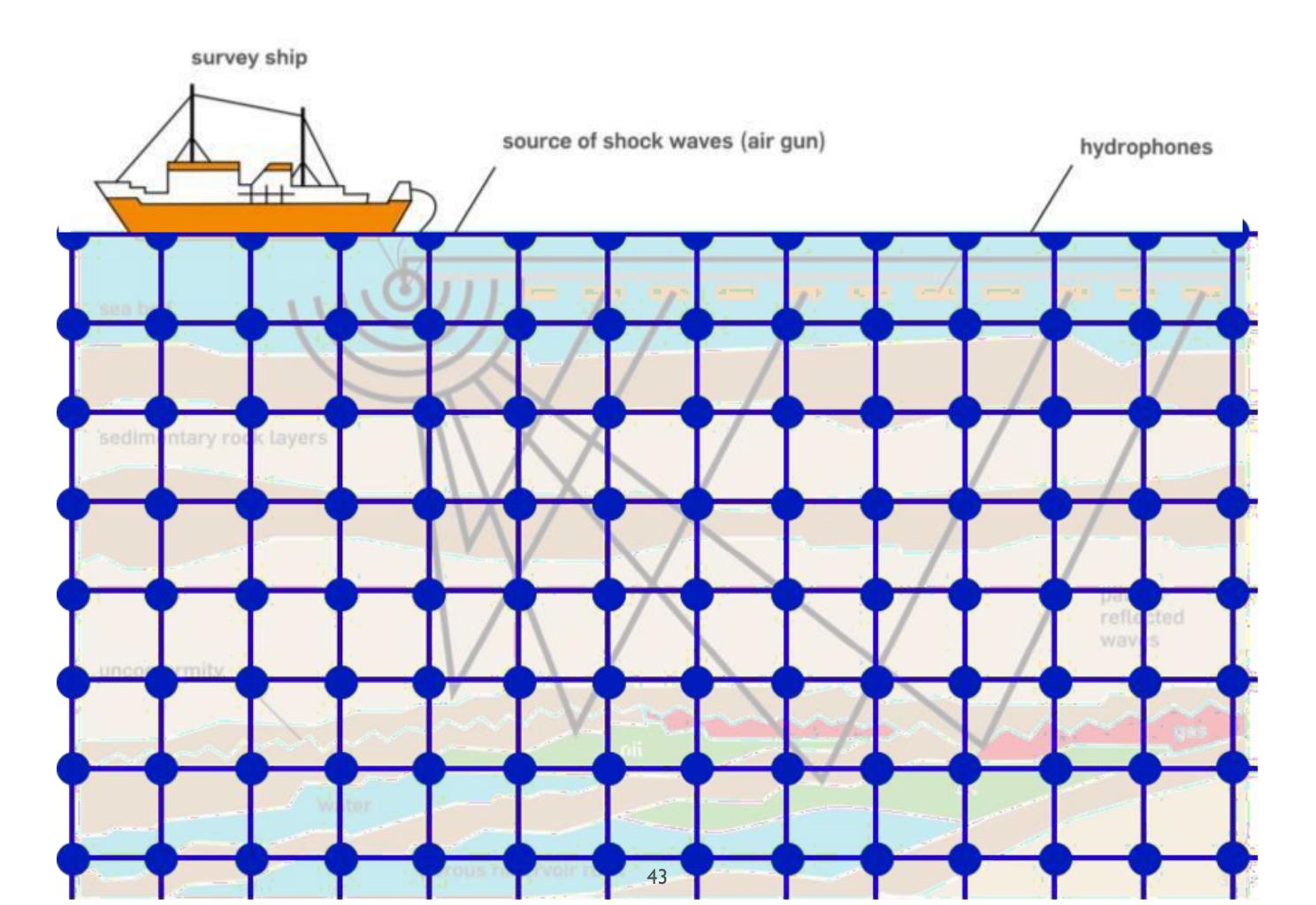
stencil_2 = Eq(usol.forward, solve(pde_2, usol.forward))

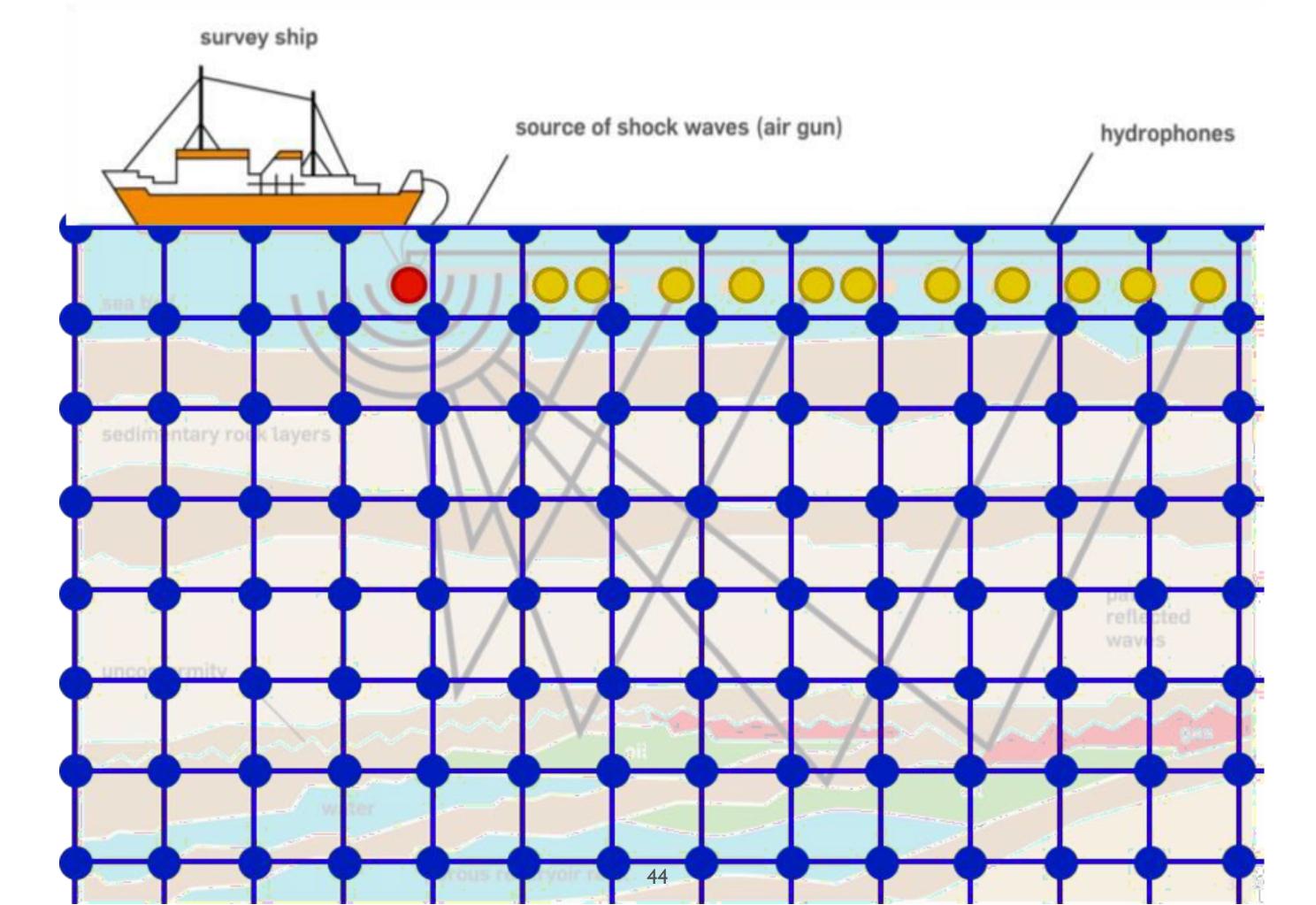
Fuse iteration spaces

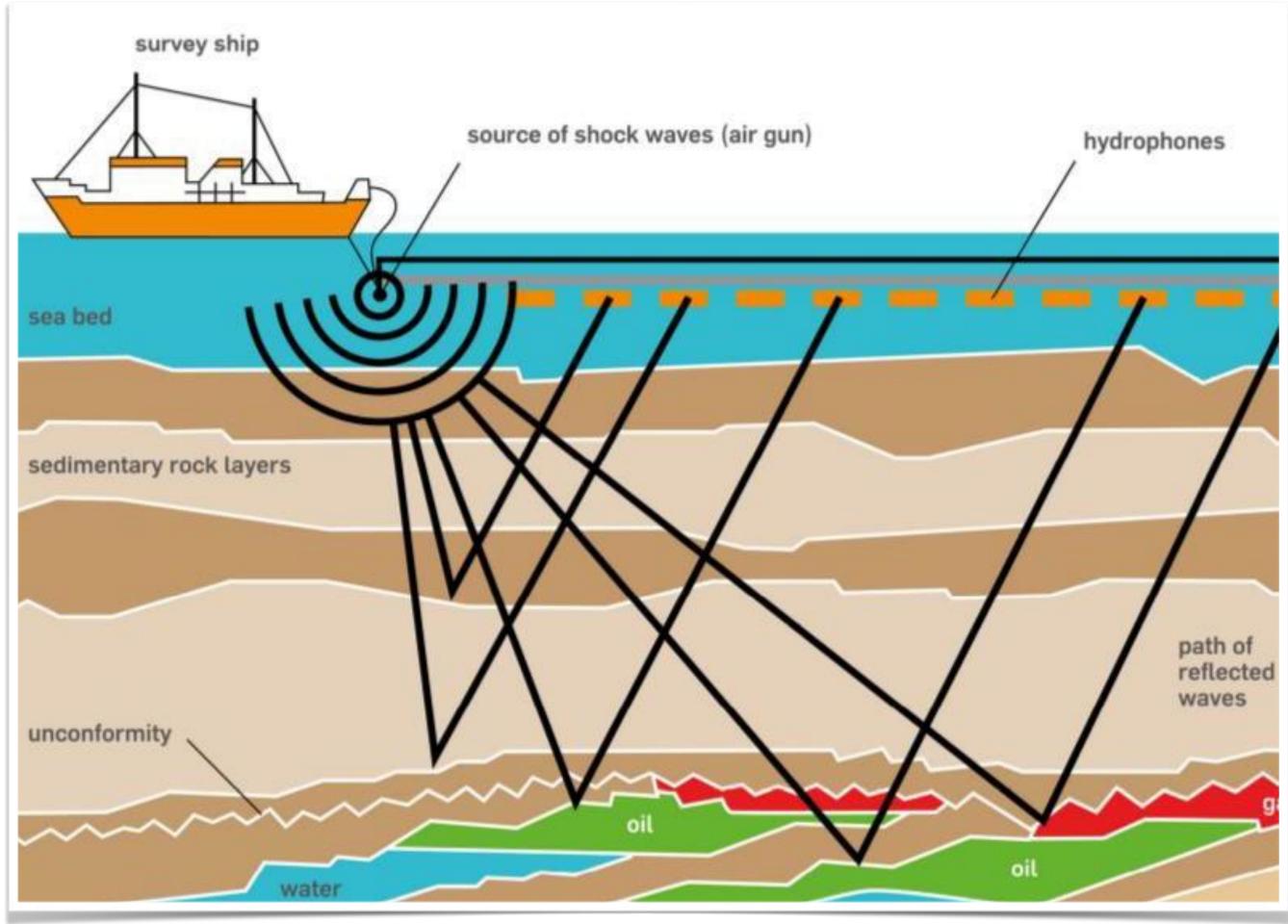
- Indirection mapping has changed. We still use indirections but now they are on the point.
- By using the aligned structure, we fuse the source injection loop inside the kernel update iteration space.
- The source mask SM is used to add (if 1) or not (if 0) the impact and SID is used to indirect to the impact values using the traversed grid coordinates.

Listing 4: Stencil kernel update with fused source injection.

1 for
$$t = 1$$
 to nt do
2 | for $x = 1$ to nx do
3 | for $y = 1$ to ny do
4 | for $z = 1$ to nz do
5 | | $A(t, x, y, z, s);$
6 | | $A(t, x, y, z, s);$
6 | | $Ior z^2 = 1$ to nz do
7 | | | $u[t, x, y, z^2] + = SM[x, y, z^2] * src_dcmp[t, SID[x, y, z^2]];$







Stencils are everywhere

- Vast literature on optimizing stencils...
- Parallelism, cache optimizations, accelerators
- From simplistic (1d-3pt)
 ...to wide and
 ...complex

survey ship

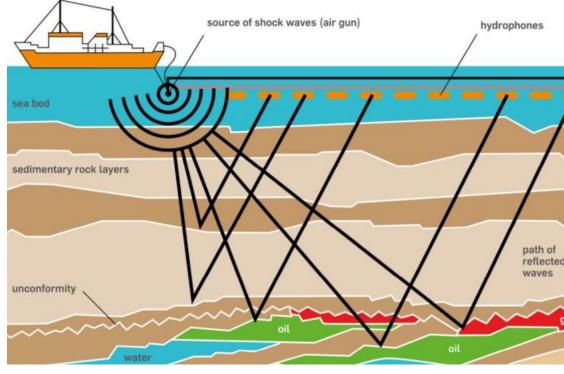


Illustration of an Offshore Seismic Survey, Source: Kris Energy

